



# Adventures & Explorations: LC

# Guideline

- <https://www.teamblind.com/post/New-Year-Gift---Curated-List-of-Top-75-LeetCode-Questions-to-Save-Your-Time-OaM1orEU>

- Steps:

- Keep calm and patient
- Read question carefully and understand requirements fully: 5-10 mins
- Set out strategy and get confirmation
- Code main logic and handle edge cases (don't hope for luck)

- Good advice/watching: <https://www.youtube.com/c/NeetCode>   <https://neetcode.io/>

- udemy course:

<https://www.udemy.com/course/datastructurescncpp/?referralCode=BD2EF8E61A98AB5E011D>

# Key steps:

- #1: Read Q clearly and understand requirements exactly, ask for clarification as needed, figure which category and which what data structure.
- #2: Discuss and confirm strategy
- #3: Code main logic first, and try to cover edge cases
- #4: don't hope for luck that it works magically, make sure/ understand every step clearly.
- #5: Correct/Improve as it runs test cases

# C++ STL: unordered\_map

```
#include <unordered_map>
unordered_map<string, int> umap;
unordered_map<string, int>::iterator itr;
pair<string, int> elem;
umap["test1"] = 1; umap["test2"] = 2;
elem = make_pair("test3", 199); umap.insert(elem); // add

itr = umap.find("mykey"); if (itr != umap.end()) printf("found mykey"); //look up

for (itr =umap.begin(); itr != umap.end(); itr++) { // iterate thru for (auto kv: umap) kv.first , kv.second
    printf("key is %s , val is %d \n", itr->first, itr->second);
umap.erase("my key"); umap.erase(umap.begin());
Umap.count("my key"); return zero if not found;
umap.size(); // return number of elements in the map
Umap.empty() ;//tell if there is anything
```

# C++ STL unordered\_set

```
#include <unordered_set>
unordered_set<string> set1;
unordered_set<string>::iterator itr1;
unordered_set<int> set2;
set1.insert("hello"); set1.insert("world"); // add
if (set1.find("myname") != set1.end()) { // found it ... } // look up
for (auto itr = set1.begin(); itr != set1.end(); itr++) { //iterate
    printf("%s", *itr);
}
set2.insert(10); set2.insert(20); set2.erase(20); set2.erase(set2.find(10));
set2.count(key) => 0 or 1; set2.size() how many elements?
if (set1.empty()) { // nothing inside ...}
```

# C++ STL Stack and Queue

```
#include <stack>
```

```
stack<int> st;
```

```
st.push(10); st.push(20); st.push(30);
```

```
while (!st.empty()) { printf("%d", st.top()); st.pop();}
```

```
st.size();
```

```
#include <queue>
```

```
queue<int> q;
```

```
q.push(100); q.push(200); q.push(300);
```

```
while (!q.empty()) { printf("%d, ", q.front()); q.pop();}
```

```
q.size();
```

# C++ STL Vector: used as ArrayList in Java

```
#include <vector>
```

```
vector<int> g1;
```

```
for (int i = 1; i <= 10; i++) g1.push_back(i * 10);
```

```
cout << "\n Reference operator [g] : g1[2] = " << g1[2];
```

```
cout << "\n Using at : g1.at(4) = " << g1.at(4);
```

```
cout << "\nfront() : g1.front() = " << g1.front();
```

```
cout << "\nback() : g1.back() = " << g1.back();
```

```
int* pos = g1.data(); // pointer to the first element
```

```
g1.push_back(15); g1.pop_back();
```

```
g1.insert(g1.begin(), 100); g1.insert(g1.begin()+3, 300);
```

```
g1.erase(g1.begin() +3) ; //erase
```

```
int index = find(g1.begin(), g1.end(), key) - g1.begin();
```

```
#include <bits/stdc++.h>
```

```
vector<int> v1 { 1, 20, 3, 40, 5, 60};
```

```
sort(v1.begin(), v1.end(), less<int>()); // less is default { 1,3,5,20, 40, 60}
```

# C++ STL string class

```
#include <string>
std::string s = "Hello";
std::string greet = s + " World"; //concatenation easy!
str.push_back('s'); str.length(); str.at(i) is same as str[i]
to_string(123); // convert integer to string
l = stoi("123"); // return integer 123;
str.find(subStr,0) == 0 => str is started with subStr
str.push_back('c');
string::append (size_type num, char c)

str.substr(pos, len);
const char *cp = str.c_str();
```



# C++ STD priority\_queue: min heap, adjust on the fly based on new value pq.push(), pq.pop()

- #include <queue>
- priority\_queue<int, vector<int>, Cmp> pq;  
// without Cmp, big num# first pop(), Cmp=less<int>, called Max heap
- Comparator:  
class Cmp {  
public:  
 bool operator()(int i1, int i2) {  
 return i1 > i2 ; //pq.top(): will be smallest number, min heap  
 }  
}
- priority\_queue<int, vector<int>, greater<int>> min\_hep;

# Binary search in vector

```
#include <iostream>
#include <vector>
#include <algorithm>
#include <iostream>

typedef std::vector<int>::iterator iter;

int main() {

    std::vector<int> vec = {10, 20, 30, 30, 20, 10, 10, 20};

    // sort the data
    // the data will be:
    // 10, 10, 10, 20, 20, 20, 30, 30
    std::sort(vec.begin(), vec.end());

    // index of the first element, greater than or equal to 20
    iter low = std::lower_bound(vec.begin(), vec.end(), 20);

    // index of the first element, greater than 20
    iter high = std::upper_bound(vec.begin(), vec.end(), 20);

    std::cout << "index of first element, greater than or equal to 20 is: " << (low - vec.begin()) << '\n';

    std::cout << "index of first element, greater than to 20 is: " << (high - vec.begin()) << '\n';

    // classic binary search
    // check whether a given value exists in the array or not
    if (std::binary_search(vec.begin(), vec.end(), 99)) {
        std::cout << "Found\n";
    } else {
        std::cout << "Not found\n";
    }
}
```

- sort: you can use binary search only on a *sorted* data, so you must guarantee that the data is sorted, before searching.
- lower\_bound: this function returns an *iterator* to the first element that is **greater than or equal to** value.
- upper\_bound: this function returns an *iterator* to the first element that is **greater than** value.
- binary\_search:: this function returns a *boolean*, whether the value is found or not (exactly as your program).

Return 3

Return 6

# Strings

```
char str[] = "HELLO";
```



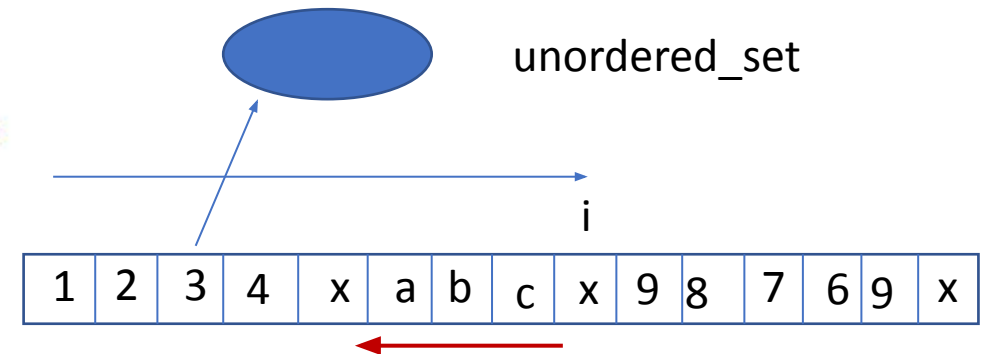
Beginning Of String

End Of String

# LC String#1: 3. Longest Substring Without Repeating Characters

```
1 class Solution {
2 public:
3     int lengthOfLongestSubstring(string s) {
4         int res = 0;
5         int c = 0;
6         if (s.length() < 2) return s.length();
7         unordered_set<char> seen;
8         for (int i = 0; i < s.length(); i++) {
9             if (seen.find(s.at(i)) == seen.end()) {
10                 c++;
11                 res = max(c, res);
12                 seen.insert(s.at(i));
13             } else {
14                 seen.clear();
15                 // search back for the last occurrence of this character
16                 seen.insert(s.at(i));
17                 c = 1;
18                 for (int j = i-1; j >= 0; j--) {
19                     if (s.at(j) != s.at(i)) {
20                         seen.insert(s.at(j));
21                         c++;
22                     } else {
23                         break;
24                     }
25                 }
26             }
27         }
28         return res;
29     }
30 };
```

- Scan thru string, push each one into “set” named as “**seen**” if not seen, c++
- If it is seen before, reset count “c”, and clear “seen”, **search back** until current char is hit(**where is i is adjusted to**) add all back into “seen”



# LC String#2: two pointers/ Caterpillar algorithm

## 424. Longest Repeating Character with K Replacement

You are given a string `s` and an integer `k`. You can choose any character of the string and change it to any other uppercase English character. You can perform this operation at most `k` times.

Return the length of the longest substring containing the same letter you can get after performing the above operations.

### Example 1:

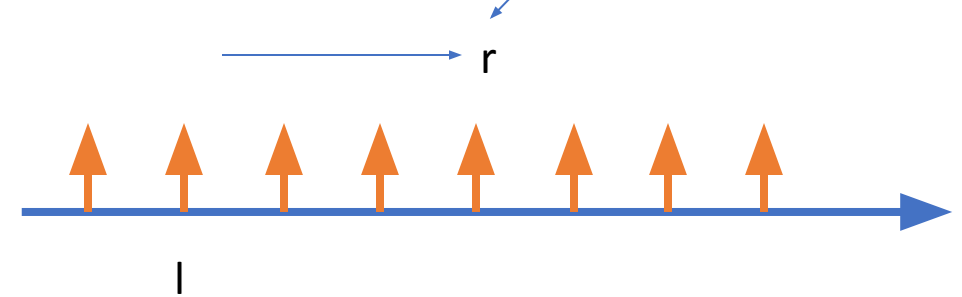
Input: `s = "ABAB", k = 2`

Output: 4

Explanation: Replace the two 'A's with two 'B's or vice versa.

```
1  class Solution {
2  public:
3      int characterReplacement(string s, int k) {
4          int L = s.length();
5          int res = 0;
6          int l = 0, r = 0;
7          int counters[26] = { 0 };
8          bool adv_right = true;
9          while (l <= r && r < L) {
10             if (adv_right)
11                 counters[s.at(r) - 'A']++;
12             else
13                 counters[s.at(l-1) - 'A']--;
14
15             //find max freq of current window
16             int max_freq = 0;
17             for (int i = 0; i < 26; i++) {
18                 max_freq = max(max_freq, counters[i]);
19             }
20             int rc = r-l+1 - max_freq;
21             if (rc <= k) {
22                 res = max(r-l+1, res);
23                 printf("\n res=%d, l=%d, r=%d", res, l, r);
24                 adv_right = true;
25                 r++;
26             } else {
27                 printf("\n advance left l=%d", l+1);
28                 l++;
29                 adv_right = false;
30             }
31         }
32         return res;
33     }
34 }
```

- Use Sliding windows:  $l=0, r=0 \rightarrow L$
- Use a counters[26] array to count each **inside window's letters** frequency when  $l/r$  is advanced
- Inside a sliding window, find the `max_freq` letter and its counter, window width ( $w=r-l+1$ ),  **$w - \text{max\_freq} \leq k$ , advance  $r$** , other wise advance  $l$





# LC string#3: two pointers/ Caterpillar algorithm

## 76. Minimum Window Substring: from S including T

```
1  class Solution {
2  public:
3      string minWindow(string s, string t) {
4          int l = 0, r = 0, res = INT_MAX, r_ = 0, l_ = 0;
5          bool adv_r = true;
6          int LEN = 'z' - 'A' + 1;
7          int t_counters['z' - 'A' + 1] = { 0 };
8          int s_counters['z' - 'A' + 1] = { 0 };
9
10         for (int i = 0; i < t.length(); i++) t_counters[t.at(i) - 'A']++;
11         while (l <= r && r < s.length()) {
12             if (adv_r)
13                 s_counters[s.at(r) - 'A']++;
14             else
15                 s_counters[s.at(l-1) - 'A']--;
16
17             //check if current window has all string t
18             bool cover_t = true;
19             for (int i = 0; i < LEN; i++) {
20                 if (t_counters[i] != 0 && s_counters[i] < t_counters[i]) {
21                     cover_t = false;
22                     break;
23                 }
24             }
25             if (cover_t) {
26                 if (r-l+1 < res) {
27                     r_ = r; l_ = l;
28                     res = r - l + 1;
29                 }
30                 adv_r = false;
31                 l++;
32             } else {
33                 adv_r = true;
34                 r++;
35             }
36         }
37
38         return res == INT_MAX? "":s.substr(l_, res);
39     }
40 }
```

Given two strings `s` and `t` of lengths `m` and `n` respectively, return the **minimum window substring** of `s` such that every character in `t` (**including duplicates**) is included in the window. If there is no such substring, return the empty string `""`.

The testcases will be generated such that the answer is **unique**.

A **substring** is a contiguous sequence of characters within the string.

### Example 1:

Input: `s = "ADOBECODEBANC"`, `t = "ABC"`

Output: `"BANC"`

Explanation: The minimum window substring `"BANC"` includes 'A', 'B', and 'C' from string `t`.

- **Sliding window**: expand R to cover T, write down W
- Advance L until not cover T, write down smaller W,
- Go back to first step until reach the end of string.

## LC string#4: 242. Valid Anagram: s = "anagram", t = "nagaram" => true

```
1  class Solution {  
2  public:  
3      bool isAnagram(string s, string t) {  
4          if (s.length() != t.length()) return false;  
5          int s_c['z'-'a'+1] = {0};  
6          int t_c['z'-'a'+1] = {0};  
7          for (int i = 0; i < s.length(); i++) {  
8              s_c[s.at(i)-'a']++;  
9              t_c[t.at(i)-'a']++;  
10         }  
11         return memcmp(s_c, t_c, sizeof(int)*('z'-'a'+1)) == 0;  
12     }  
13 };
```

An Anagram is a word or phrase formed by rearranging the letters of a different word or phrase, typically using all the original letters exactly once.

# LC string#5: 49. Group Anagrams:

Input: strs = ["eat","tea","tan","ate","nat","bat"]  
Output: [["bat"],["nat","tan"],["ate","eat","tea"]]

```
1  class Solution {  
2  public:  
3      vector<vector<string>> groupAnagrams(vector<string>& strs) {  
4          vector<vector<string>> res;  
5          vector<vector<int>> counters;  
6          vector<bool> visited;  
7          for (int i = 0; i < strs.size(); i++) {  
8              string s = strs[i];  
9              vector<int> c(26,0);  
10             for (int j = 0; j < s.length(); j++)  
11                 c[s.at(j) - 'a']++;  
12             counters.push_back(c);  
13             visited.push_back(false);  
14         }  
15  
16         for (int i = 0; i < strs.size(); i++) {  
17             if (!visited[i]) {  
18                 vector<string> ans;  
19                 ans.push_back(strs[i]);  
20                 visited[i] = true;  
21                 for (int j = i+1; j < strs.size(); j++) {  
22                     if (counters[i] == counters[j]) {  
23                         ans.push_back(strs[j]);  
24                         visited[j] = true;  
25                     }  
26                 }  
27                 res.push_back(ans);  
28             }  
29         }  
30         return res;  
31     }  
32 };
```

- **Character counters**: for each string **vector<int> c(26,0)**
- Visited[] Boolean to speed up
- Scan thru strs and counters[]  
**vector compare**: counters[i] == counters[j]



# LC string#6: 20. Valid Parentheses:

Input: s = "()[]{}"  
Output: true

Input: s = "([)]"  
Output: false

```
1 class Solution {  
2 public:  
3     bool isValid(string s) {  
4         stack<char> st;  
5  
6         for (int i = 0; i < s.length(); i++) {  
7             if (s.at(i) == '(' || s.at(i) == '[' || s.at(i) == '{')  
8                 st.push(s.at(i));  
9             else {  
10                if (st.empty()) return false;  
11                char c = st.top(); st.pop();  
12                if ((c == '{' && s.at(i) == '}') ||  
13                    (c == '[' && s.at(i) == ']') ||  
14                    (c == '(' && s.at(i) == ')'))  
15                    continue;  
16                } else  
17                    return false;  
18            }  
19        }  
20        return st.empty();  
21    }  
22 }  
23 };
```

- **Use stack:** opening one, push, closing one pop and compare.

# LC string#7: 125. Valid Palindrome:

**Input:** s = "A man, a plan, a canal: Panama"

**Output:** true

**Explanation:** "amanaplanacanalpanama" is a palindrome.

```
1  class Solution {
2  public:
3      bool isAlphaNum(char c) {
4          return ((c >= 'A' && c <= 'Z') ||
5                  (c >= 'a' && c <= 'z') ||
6                  (c >= '0' && c <= '9'));
7      }
8      bool equalIgnoreCase(char a, char b) {
9          // convert to uppercase to compare
10         if (a >= 'a' && a <= 'z') {
11             a = 'A' + (a - 'a');
12         }
13         if (b >= 'a' && b <= 'z') {
14             b = 'A' + (b - 'a');
15         }
16         return a == b;
17     }
18     | bool isPalindrome(string s) {
19         int l = 0;
20         int r = s.length() - 1;
21
22         while (l < r) {
23             while (l < s.length() && !isAlphaNum(s.at(l))) l++;
24             while (r >= 0 && !isAlphaNum(s.at(r))) r--;
25             if (l >= s.length() || r < 0) break;
26             if (!equalIgnoreCase(s.at(l), s.at(r))) return false;
27             l++;
28             r--;
29         }
30         return true;
31     }
32 };
```

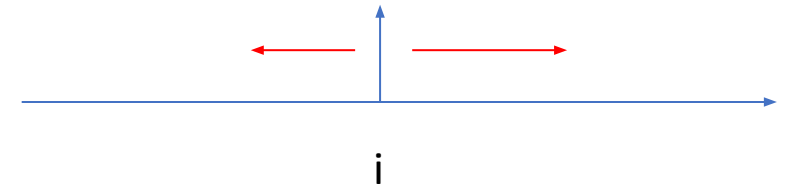
# LC string#8: 5. Longest Palindromic substring:

```
1  class Solution {  
2  public:  
3      int expandHelper(string s, int l, int r) {  
4          while (l >= 0 && r <= s.length()-1 && s.at(l) == s.at(r) ) {  
5              l--; r++;  
6          }  
7          return r-l-1;  
8      }  
9      string longestPalindrome(string s) {  
10         int len = 1;  
11         int pos = 0;  
12         for (int i = 0; i < s.length(); i++) {  
13             int len_odd = expandHelper(s, i, i);  
14             int len_even = expandHelper(s, i, i+1);  
15             int max_len = max(len_odd, len_even);  
16             if (max_len > len) {  
17                 len = max_len;  
18                 pos = i - (len-1)/2;  
19             }  
20         }  
21         return s.substr(pos, len);  
22     }  
23 };
```

Expand with l as center “odd”

Expand with i and i+1 as “even”

take longer one , record it along the way  
to get len,  $pos = i - (len-1)/2$



# LC string#9: 647: Palindromic substring:

Given a string `s`, return the number of **palindromic substrings** in it.

A string is a **palindrome** when it reads the same backward as forward.

A **substring** is a contiguous sequence of characters within the string.

## Example 1:

Input: `s = "abc"`

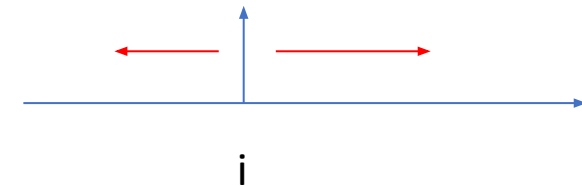
Output: 3

Explanation: Three palindromic strings: "a", "b", "c".

```
1 class Solution {
2 public:
3     void expandHelper(string s, int l, int r, int & result) {
4         while ( l >=0 && r < s.length() && s.at(l) == s.at(r)) {
5             result++;
6             l--; r++;
7         }
8     }
9     int countSubstrings(string s) {
10        int res = 0;
11
12        for (int i = 0; i < s.length(); i++) {
13            expandHelper(s, i, i, res);
14            expandHelper(s, i, i+1, res);
15        }
16
17        return res;
18    }
19};
```

Expand with `i` as center “odd”

Expand with `i` and `i+1` as “even”



# LC string#10: 271: Encode and Decode strings:

```
1  class Codec {
2  public:
3
4      // Encodes a list of strings to a single string.
5      string encode(vector<string>& strs) {
6          string res;
7          for (int i = 0; i < strs.size(); i++) {
8              string tmp = to_string(strs[i].length())+"@"+ strs[i];
9              res.append(tmp);
10         }
11         return res;
12     }
13
14     // Decodes a single string to a list of strings.
15     vector<string> decode(string s) {
16         string dec_s = s;
17
18         vector<string> res;
19         int i = 0;
20         while (i < dec_s.length()) {
21             int l = 0;
22             while (dec_s.at(i+l) != '@') {
23                 l++;
24             }
25             string num_s = dec_s.substr(i,l);
26             int len = stoi(num_s);
27             i += l+1; // skip @
28             res.push_back(dec_s.substr(i, len));
29             i += len;
30         }
31         return res;
32     }
33 };
```

Design an algorithm to encode a **list of strings** to a **string**. The encoded string is then sent over the network and is decoded back to the original list of strings.

Machine 1 (sender) has the function:

```
string encode(vector<string> strs) {
    // ... your code
    return encoded_string;
}
```

Machine 2 (receiver) has the function:

```
vector<string> decode(string s) {
    //... your code
    return strs;
}
```



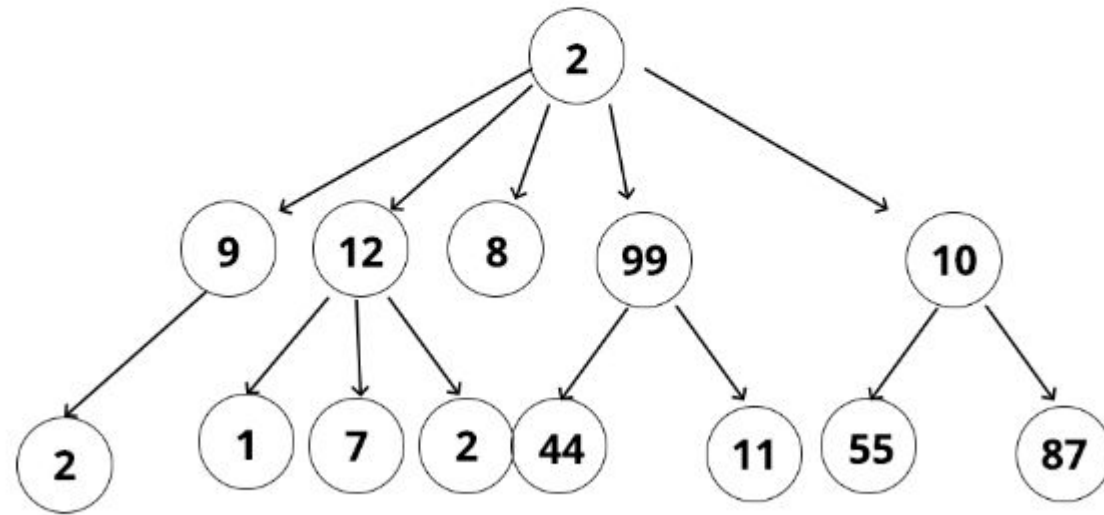
# LC string#11: 13. Roman to Integer:

```
1 class Solution {
2 public:
3     int romanToInt(string s) {
4         unordered_map<string, int> map;
5
6         map["I"] = 1; map["V"] = 5; map["IV"] = 4; map["IX"] = 9;
7         map["X"] = 10; map["L"] = 50; map["XL"] = 40; map["XC"] = 90;
8         map["C"] = 100; map["D"] = 500; map["CD"] = 400; map["CM"] = 900;
9         map["M"] = 1000;
10        int i = 0;
11        int res = 0;
12        while (i < s.length()) {
13            // check if it one or two characters
14            if((i+1) < s.length()) {
15                string dStr = s.substr(i,2);
16                if (map.find(dStr) != map.end()) {
17                    res += map[dStr];
18                    i += 2;
19                    continue;
20                }
21            }
22            res += map[s.substr(i,1)];
23            i++;
24        }
25        return res;
26    }
27};
```

# 12. Integer to Roman

```
1 class Solution {
2 public:
3     string intToRoman(int num) {
4         vector<int> values { 1000, 900, 500, 400, 100,
5                             90, 50, 40, 10, 9, 5, 4, 1};
6         vector<string> syms { "M", "CM", "D", "CD", "C", "XC",
7                               "L", "XL", "X", "IX", "V", "IV", "I"};
8         string res = "";
9         for (int i = 0; i < values.size() && num > 0; i++) {
10             while (num >= values[i]) {
11                 num -= values[i];
12                 res.append(syms[i]);
13             }
14         }
15         return res;
16     }
17};
```

# TREES

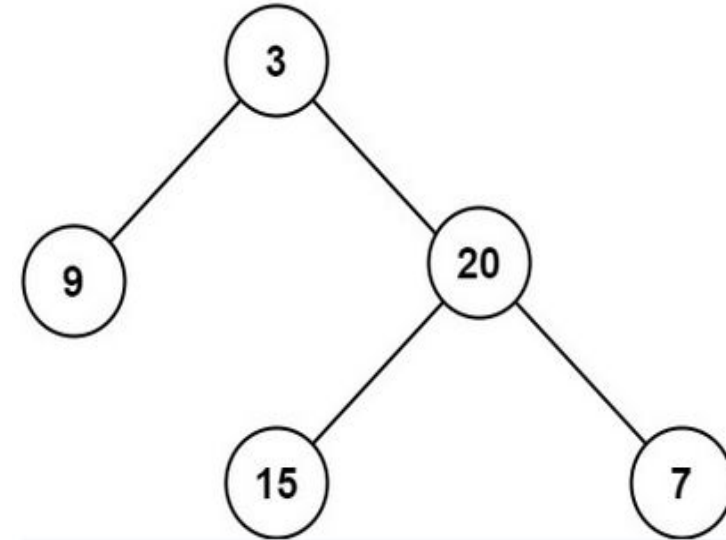


# LC tree#1: 104. Maximum Depth of Binary Tree:

```
12 ▾ class Solution {  
13     public:  
14 ▾     int maxDepth(TreeNode* root) {  
15         int max_level = 0;  
16  
17         if (!root) return max_level;  
18         queue<TreeNode*> Q;  
19         Q.push(root);  
20  
21 ▾         while (!Q.empty()) {  
22             max_level++;  
23             int s = Q.size();  
24 ▾             for (int i = 0; i < s; i++) {  
25                 TreeNode *n = Q.front(); Q.pop();  
26                 if (n->left) Q.push(n->left);  
27                 if (n->right) Q.push(n->right);  
28             }  
29         }  
30  
31         return max_level;  
32     }  
33 };
```

Use Q to do a level traversal and  
Counting levels

Example 1:



Input: root = [3,9,20,null,null,15,7]

Output: 3

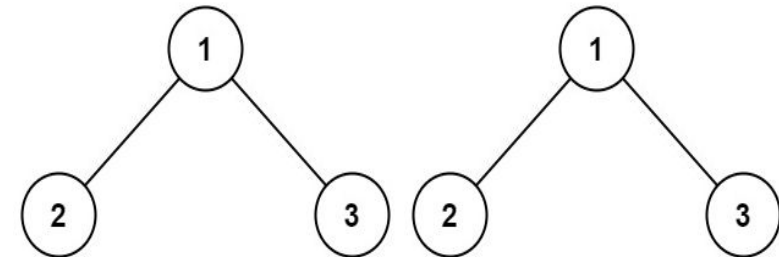


# LC tree#2: 100. Same Tree:

```
12 class Solution {
13 public:
14     void dfs(TreeNode *root, vector<int>& v) {
15         if (root == NULL) return;
16         TreeNode *dummy = new TreeNode(INT_MAX);
17
18         stack<TreeNode*> st;
19         st.push(root);
20         while (!st.empty()) {
21             TreeNode *n = st.top(); st.pop();
22             v.push_back(n->val);
23             if (n->val != INT_MAX) {
24                 st.push(n->left? n->left:dummy);
25                 st.push(n->right? n->right:dummy);
26             }
27         }
28         delete(dummy);
29     }
30     bool isSameTree(TreeNode* p, TreeNode* q) {
31         vector<int> p_vals, q_vals;
32         dfs(p, p_vals);
33         dfs(q, q_vals);
34         return p_vals == q_vals;
35     }
36 };
```

Do a **DFS using stack** for each tree, **for absent node**, **Add a dummy node with INT\_MAX as flag**, which should do nothing except add INT\_MAX into vector after it is popped out from stack

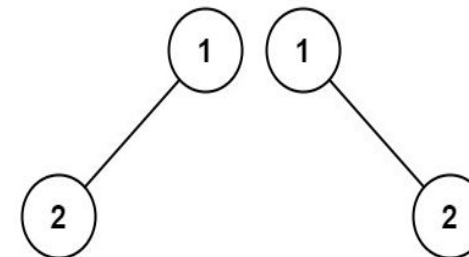
Example 1:



Input: p = [1,2,3], q = [1,2,3]

Output: true

Example 2:



Input: p = [1,2], q = [1,null,2]

Output: false

# LC tree#3: 226. Invert Binary Tree:

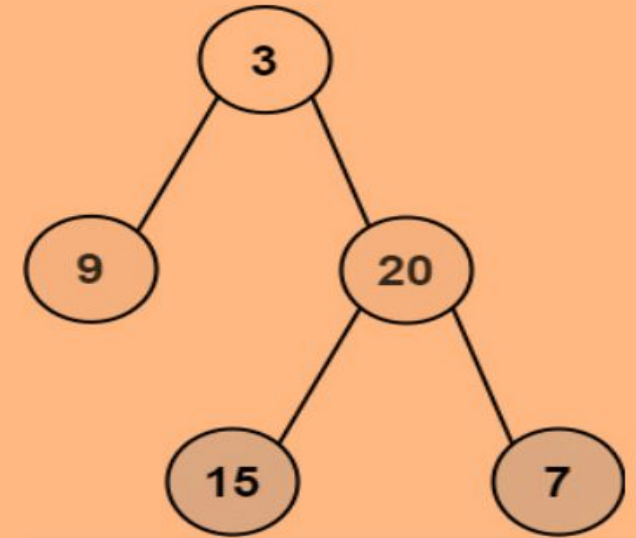
```
12 ▼ class Solution {  
13     public:  
14 ▼     TreeNode* invertTree(TreeNode* root) {  
15         if (root == NULL) return root;  
16         TreeNode *t = root->left;  
17         root->left = root->right;  
18         root->right = t;  
19         invertTree(root->left);  
20         invertTree(root->right);  
21     }  
22     return root;  
23 }  
24 };
```

Exchange left with right node  
Do recursion for left and right.

# LC tree#4.1: 102. Binary Tree Level Order Traversal:

```
12 ▾ class Solution {  
13     public:  
14 ▾     vector<vector<int>> levelOrder(TreeNode* root) {  
15         vector<vector<int>> res;  
16         if (root == NULL) return res;  
17         queue<TreeNode*> Q;  
18         Q.push(root);  
19  
20 ▾         while (!Q.empty()) {  
21             int N = Q.size();  
22             vector<int> level;  
23 ▾             for (int i = 0; i < N; i++) {  
24                 TreeNode *n = Q.front(); Q.pop();  
25                 level.push_back(n->val);  
26                 if (n->left) Q.push(n->left);  
27                 if (n->right) Q.push(n->right);  
28             }  
29             res.push_back(level);  
30         }  
31  
32         return res;  
33     }  
34 };
```

Example 1:



**Input:** root = [3,9,20,null,null,15,7]

**Output:** [[3],[9,20],[15,7]]

Using **Q** to do level traversal  
Check **Q size** in the beginning of loop,  
which is number of nodes in that level.

# LC string#4.2: 144. Binary Tree preOrder Traversal:

```
12 ▾ class Solution {
13     public:
14 ▾         vector<int> preorderTraversal(TreeNode* root) {
15             vector<int> res;
16             if (root == NULL) return res;
17     #ifdef RECURSIVE
18             res.push_back(root->val);
19             vector<int> left = preorderTraversal(root->left);
20             for (int i = 0; i < left.size(); i++)
21                 res.push_back(left[i]);
22             vector<int> right = preorderTraversal(root->right);
23             for (int i = 0; i < right.size(); i++)
24                 res.push_back(right[i]);
25     #endif
26             TreeNode *cur = root;
27             stack<TreeNode *> st;
28 ▾             while (cur != NULL || !st.empty()) {
29 ▾                 if (cur != NULL) {
30                     res.push_back(cur->val);
31                     st.push(cur);
32                     cur = cur->left; // push down all the way to bottom most left node
33 ▾                 } else {
34                     cur = st.top(); st.pop();
35                     cur = cur->right;
36                 }
37             }
38             return res;
39         }
40     }
41 };
```

DFS using stack is cleanest!

```
42         TreeNode *cur = root;
43         stack<TreeNode *> st;
44         st.push(root);
45 ▾         while (!st.empty()) {
46             TreeNode *n = st.top(); st.pop();
47             res.push_back(n->val);
48             if (n->right) st.push(n->right);
49             if (n->left) st.push(n->left);
50         }
51         return res;
52     }
```



## LC string#4.3: 94. Binary Tree Level InOrder Traversal:

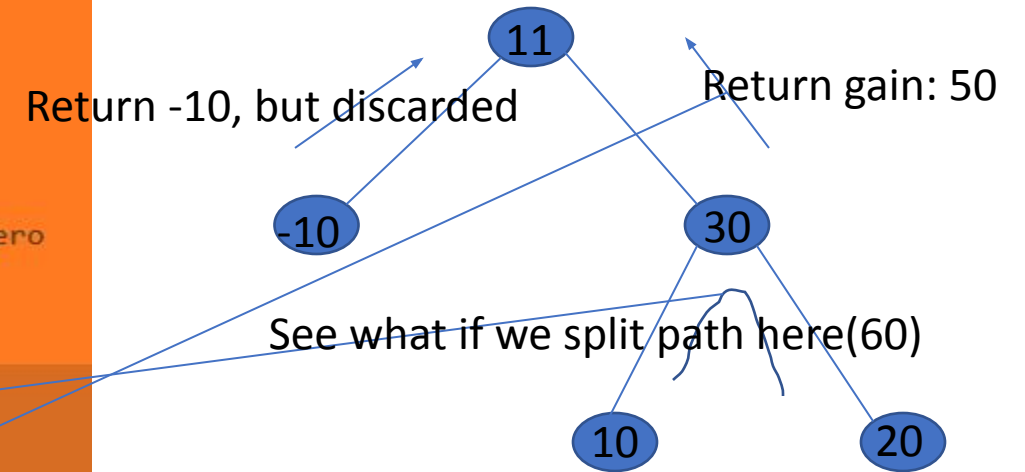
```
12 ▾ class Solution {
13     public:
14 ▾     vector<int> inorderTraversal(TreeNode* root) {
15         vector<int> res;
16         if (root == NULL) return res;
17     #ifdef RECURSIVE
18         vector<int> left = preorderTraversal(root->left);
19         for (int i = 0; i < left.size(); i++)
20             res.push_back(left[i]);
21         res.push_back(root->val);
22         vector<int> right = preorderTraversal(root->right);
23         for (int i = 0; i < right.size(); i++)
24             res.push_back(right[i]);
25     #endif
26         TreeNode *cur = root;
27         stack<TreeNode *> st;
28 ▾     while (cur != NULL || !st.empty()) {
29 ▾         if (cur != NULL) {
30             st.push(cur);
31             cur = cur->left; // push down all the way to bottom most left node
32 ▾         } else {
33             cur = st.top(); st.pop();
34             res.push_back(cur->val);
35             cur = cur->right;
36         }
37     }
38     return res;
39 }
40 }
41 };
```

# LC string#4.4: 102. Binary Tree PostOrder Traversal:

```
12 ▾ class Solution {
13     public:
14 ▾     vector<int> postorderTraversal(TreeNode* root) {
15         vector<int> res;
16         if (root == NULL) return res;
17     #ifdef RECURSIVE
18         vector<int> left = preorderTraversal(root->left);
19         for (int i = 0; i < left.size(); i++)
20             res.push_back(left[i]);
21         vector<int> right = preorderTraversal(root->right);
22         res.push_back(root->val);
23         for (int i = 0; i < right.size(); i++)
24             res.push_back(right[i]);
25     #endif
26
27         stack<TreeNode*> st;
28         st.push(root);
29 ▾         while (!st.empty()) {
30             TreeNode* cur = st.top();
31             st.pop();
32             res.insert(res.begin(), cur->val); // this makes all existing one shift to right
33             if (cur->left) st.push(cur->left); // left node first push
34             if (cur->right) st.push(cur->right);
35         }
36
37         return res;
38     }
39 };
```

# LC tree#5: 124. Binary Tree Maximum Path Sum(HARD):

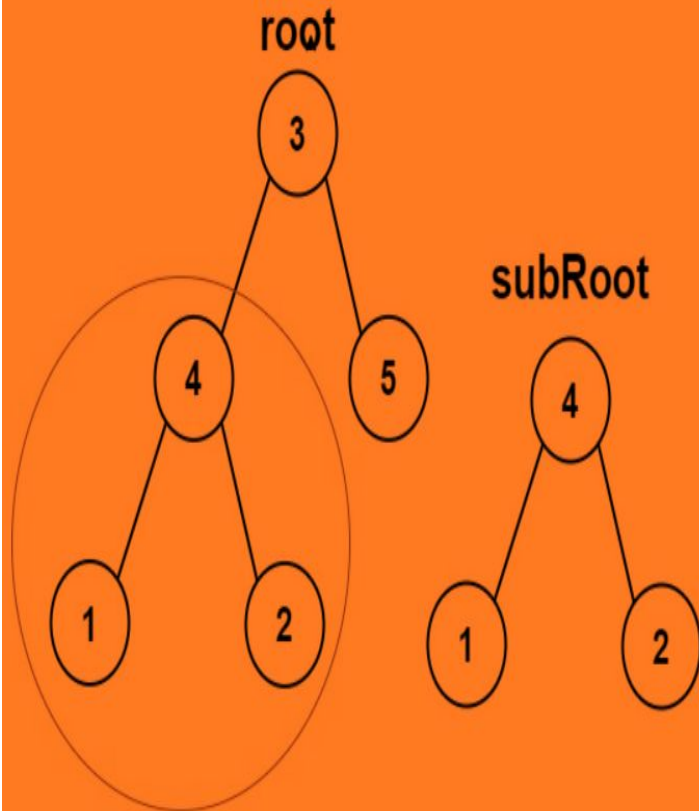
```
12 class Solution {  
13     int res = INT_MIN;  
14 public:  
15     // return gain of node without split path  
16     int max_gain(TreeNode *n) {  
17         if (n == NULL) return 0;  
18  
19         // if gain is negative, we will not take it, so set to zero  
20         int left_gain = max(max_gain(n->left), 0);  
21         int right_gain = max(max_gain(n->right), 0);  
22  
23         // check what is value if we split path at this node  
24         if (n->val + left_gain + right_gain > res)   
25             res = n->val + left_gain + right_gain;  
26  
27         return n->val + max(left_gain, right_gain);  
28     }  
29  
30     int maxPathSum(TreeNode* root) {  
31         if (root == NULL) return 0;  
32  
33         max_gain(root);  
34  
35         return res;  
36     }  
37 };
```



# LC tree#6:

## 572. Subtree of Another Tree:

Example 1:



Input: root = [3,4,5,1,2], subRoot = [4,1,2]

Output: true

- Generate list of nodes starting from root
- Using each node to do same tree checking
- Same tree check using traversal to generate vector<int> to compare.

```
12 class Solution {
13 public:
14     vector<int> levelTraver(TreeNode* root) {
15         vector<int> res;
16         if (root == NULL) return res;
17         queue<TreeNode*> Q;
18         Q.push(root);
19         TreeNode *dummy = new TreeNode(INT_MAX);
20         while (!Q.empty()) {
21             TreeNode *n = Q.front(); Q.pop();
22             res.push_back(n->val);
23             if (n->val != INT_MAX) {
24                 Q.push(n->left ? n->left:dummy);
25                 Q.push(n->right ? n->right:dummy);
26             }
27         }
28         return res;
29     }
30     bool sameTree(TreeNode *p, TreeNode *q) {
31         vector<int> p_v = levelTraver(p);
32         vector<int> q_v = levelTraver(q);
33         return p_v == q_v;
34     }
35     bool isSubtree(TreeNode* root, TreeNode* subRoot) {
36         vector<TreeNode*> nodes;
37         if (root == NULL) return false;
38
39         queue<TreeNode*> Q;
40         Q.push(root);
41         while (!Q.empty()) {
42             TreeNode *n = Q.front(); Q.pop();
43             nodes.push_back(n);
44             if (n->left) Q.push(n->left);
45             if (n->right) Q.push(n->right);
46         }
47         for (int i = 0; i < nodes.size(); i++) {
48             if (nodes[i]->val == subRoot->val) {
49                 if (sameTree(nodes[i], subRoot))
50                     return true;
51             }
52         }
53         return false;
54     }
55 }
```



# LC tree#7: 105. Construct Binary Tree from Preorder and Inorder Traversal:

Given preorder and inorder, this precisely defines a tree

First value from preorder is ROOT, which can be used

To find the partition of LEFT and RIGHT!

If only preorder/inorder given, unless it is a full balanced tree

If postorder is given instead of preorder, find root using last!

Time & Space:  $O(n)$

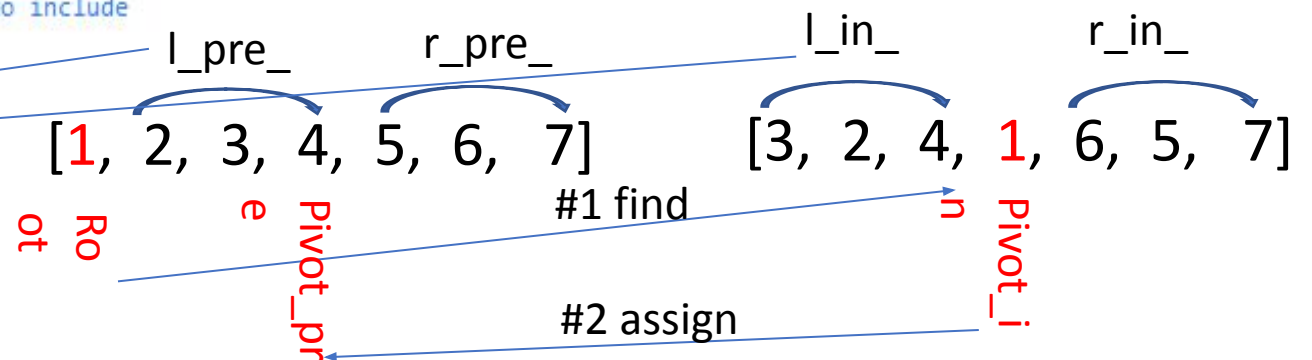
#1

```
12 class Solution {
13 public:
14     TreeNode* buildTree(vector<int>& preorder, vector<int>& inorder) {
15         //first or preorder is for root
16         if (preorder.size() == 0) return NULL;
17
18         TreeNode *root = new TreeNode(preorder[0]); if (preorder.size() == 1) return root;
19         // find the pivot position of root node inside inorder
20         // (because of unique value) using root value
21         int pivot_in = find(inorder.begin(), inorder.end(), preorder[0]) - inorder.begin();
22         //partition preorder, use pivot_in (inclusive)
23         int pivot_pre = pivot_in;
24         vector<int> l_pre_, r_pre_;
25         for (int j = 1; j <= pivot_pre; j++) // exclude root, start with 1
26             l_pre_.push_back(preorder[j]);
27         for (int j = pivot_pre+1; j < preorder.size(); j++)
28             r_pre_.push_back(preorder[j]);
29         //partition inorder, pivot is pivot_in
30         vector<int> l_in_, r_in_;
31         for (int j = 0; j < pivot_in; j++) // pivot_in is root, not to include
32             l_in_.push_back(inorder[j]);
33         for (int j = pivot_in+1; j < inorder.size(); j++)
34             r_in_.push_back(inorder[j]);
35
36         root->left = buildTree(l_pre_, l_in_);
37         root->right = buildTree(r_pre_, r_in_);
38
39         return root;
40     }
41 }
```

#2

preorder

inorder



# LC tree#8: 226. serialize and deserialize binary tree:

```
10 class Codec {
11 public:
12
13     // Encodes a tree to a single string following preorder sequence
14     string serialize(TreeNode* root) {
15         if (root == NULL) return "9999,"; // flag as end
16
17         string res = to_string(root->val) + ",";
18         res.append(serialize(root->left));
19         res.append(serialize(root->right));
20         return res;
21     }
22     int pos = 0; // use to track current processed node
23     TreeNode* buildTree(vector<int>& vals) {
24         if (vals[pos] == 9999) {
25             pos++; return NULL;
26         }
27         TreeNode* root = new TreeNode(vals[pos]); pos++;
28         root->left = buildTree(vals);
29         root->right = buildTree(vals);
30         return root;
31     }
32     // Decodes your encoded data to tree.
33     TreeNode* deserialize(string data) {
34         vector<int> vals;
35         int s = 0, len = 0;
36         for (int i = 0; i < data.size(); i++) {
37             if (data.at(i) != ',')
38                 len++;
39             else {
40                 vals.push_back(stoi(data.substr(s, len)));
41                 s = i+1;
42                 len = 0;
43             }
44         }
45         return buildTree(vals);
46     }
47 };
```

Use simple recursion/dfs to do preorder traversal

Using 9999 as NULL node.

Processing serialized string into vector<int> and use

Recursion to deserialize tree node.

Use variable “pos” to track position.

# LC tree#9: 98. Validate Binary Search Tree:

```
12 ▾ class Solution {  
13     public:  
14         bool helper (TreeNode *root, int64_t low, int64_t high) {  
15             if (root->val > low && root->val < high) {  
16                 if (root->left)  
17                     if (!helper(root->left, low, root->val)) return false;  
18                 if (root->right)  
19                     if (!helper(root->right, root->val, high)) return false;  
20             } else {  
21                 return false;  
22             }  
23             return true;  
24         }  
25         bool isValidBST(TreeNode* root) {  
26             if (root == NULL) return false;  
27             int64_t LOW_INF = ((int64_t)INT_MIN)-1;  
28             int64_t HIGH_INF = ((int64_t)INT_MAX) + 1;  
29             return helper(root, LOW_INF, HIGH_INF);  
30         }  
31     };
```

- Use int64\_t LOW\_INF = (int64\_t)INT\_MIN-1
- Use int64\_t HIGH\_INF = (int64\_t)INT\_MAX+1
- To check left node, need pass down current node value as HIGH
- To check right node, need pass down current node value as LOW
- If there is any false return, return back all the way.

# LC tree#10: 230. Kth Smallest Element in a BST:

```
12 class Solution {
13 public:
14
15 void inorderT(TreeNode * root, vector<int>& list) {
16     if (!root) return;
17     if (root->left) inorderT(root->left, list);
18     list.push_back(root->val);
19     if (root->right) inorderT(root->right, list);
20 }
21 int kthSmallest(TreeNode* root, int k) {
22     vector<int> list;
23     inorderT(root, list);
24     return list[k-1];
25 }
26 };
```

BST inorder traversal will give list of Value in ascending order.

```
int hight_dfs(Node* node, int &res) {
    if (node == NULL) return -1;
    if (node->children.size() == 0) return 0;

    vector<int> heights;
    for (auto n: node->children) {
        heights.push_back(hight_dfs(n, res));
    }
    sort(heights.begin(), heights.end(), greater<int>());
    if (heights.size() >= 2) {
        res = max(res, heights[0] + heights[1] + 2); // add two: one for each side
    }
    return heights[0]+1;
}
```

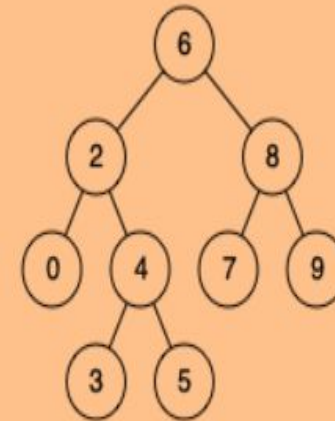


# LC tree#11: 235. Lowest Common Ancestor of a Binary Search Tree:

- Find p in BST into vector<int>
- Find q in BST into unordered\_map<int, TreeNode\*>

```
11 class Solution {
12 public:
13     void findNodeList(TreeNode * root, TreeNode *p, vector<int> &list) {
14         TreeNode *cur = root;
15         while (cur && cur->val != p->val) {
16             list.insert(list.begin(), cur->val);
17             (cur->val > p->val)? (cur = cur->left):(cur = cur->right);
18         }
19         list.insert(list.begin(), cur->val);
20     }
21     void findNodeMap(TreeNode * root, TreeNode *p, unordered_map<int, TreeNode*> &map) {
22         TreeNode *cur = root;
23         while (cur && cur->val != p->val) {
24             map[cur->val] = cur;
25             (cur->val > p->val)? (cur = cur->left):(cur = cur->right);
26         }
27         map[cur->val] = cur;
28     }
29     TreeNode* lowestCommonAncestor(TreeNode* root, TreeNode* p, TreeNode* q) {
30         vector<int> p_list;
31         unordered_map<int, TreeNode*> q_map;
32         findNodeList(root, p, p_list);
33         findNodeMap(root, q, q_map);
34         for (int i = 0; i < p_list.size(); i++){
35             if (q_map.find(p_list[i]) != q_map.end())
36                 return q_map[p_list[i]];
37         }
38         return NULL;
39     }
40 };
```

Example 1:



**Input:** root = [6,2,8,0,4,7,9,null,null,3,5], p = 2, q = 8

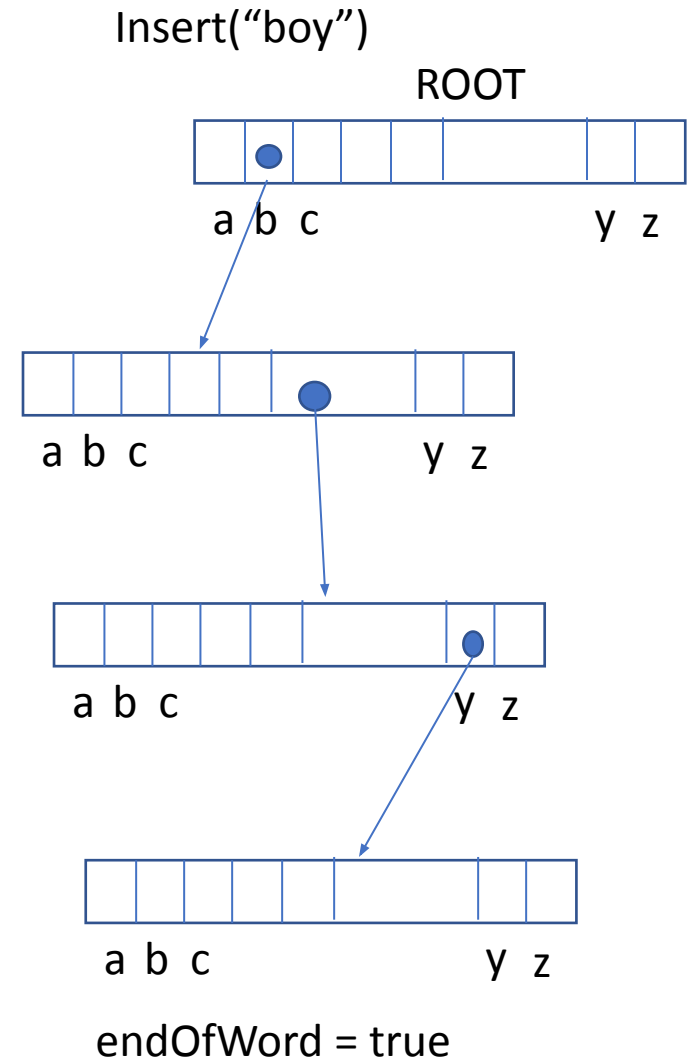
**Output:** 6

**Explanation:** The LCA of nodes 2 and 8 is 6.

236. Lowest Common Ancestor of a Binary Tree

# LC tree#12: 208. Implement Trie (Prefix Tree)

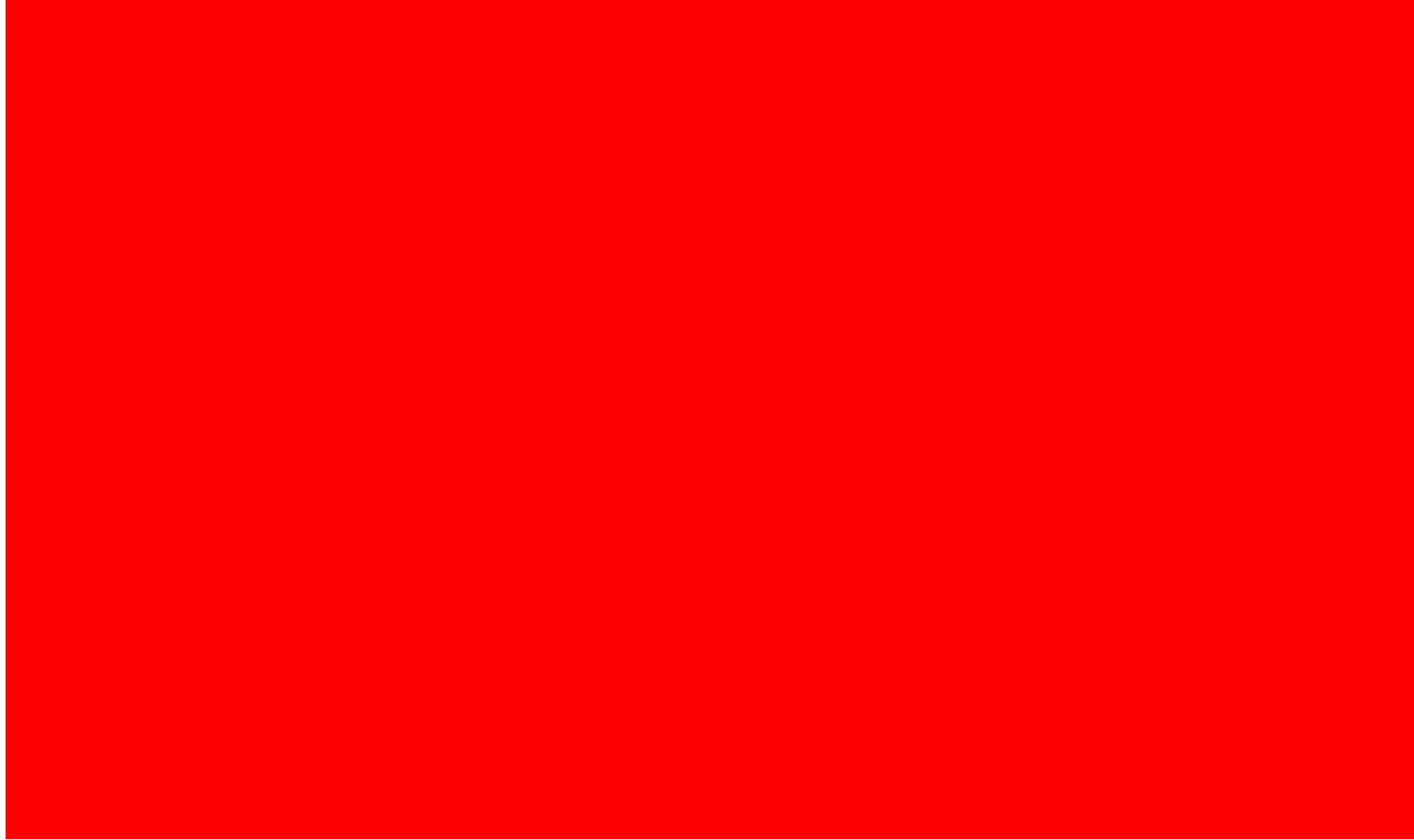
```
1 class Trie {
2     Trie* children[26];
3     bool endOfWord;
4
5 public:
6     /** Initialize your data structure here. */
7     Trie() {
8         for (int i = 0; i < 26; i++) children[i] = NULL;
9         endOfWord = false;
10    }
11    /** Inserts a word into the trie. */
12    void insert(string word) {
13        Trie *p = this;
14        for (int i = 0; i < word.length(); i++) {
15            char c = word.at(i); // each character is used as key into children
16            if (p->children[c-'a'] == NULL) {
17                Trie *t = new Trie();
18                p->children[c-'a'] = t;
19            }
20            p = p->children[c-'a'];
21        }
22        p->endOfWord = true;
23    }
24    /** Returns if the word is in the trie. */
25    bool search(string word) {
26        Trie *p = this;
27        for (int i = 0; i < word.length(); i++) {
28            char c = word.at(i);
29            p = p->children[c-'a'];
30            if (!p) return false;
31        }
32        return p->endOfWord;
33    }
34    /** Returns if there is any word in the trie that starts with the given prefix. */
35    bool startsWith(string prefix) {
36        Trie *p = this;
37        for (int i = 0; i < prefix.length(); i++) {
38            char c = prefix.at(i);
39            p = p->children[c-'a'];
40            if (!p) return false;
41        }
42        return (p != NULL);
43    }
44 };
```



# LC tree#13: 211. Design Add and Search Words Data Structure: (essential trie)

LC tree#14: 212. Word Search II: (HARD)





# LC Array#1: 1. Two Sum (easy)

Using unordered\_map<int, int> to record  
Index and value for search: O(n)

```
1  class Solution {  
2  public:  
3      vector<int> twoSum(vector<int>& nums, int target) {  
4          unordered_map<int, int> umap;  
5          unordered_map<int, int>::iterator itr;  
6          vector<int> res;  
7  
8          for (int i = 0; i < nums.size(); i++) {  
9              int reminder = target - nums[i];  
10             itr = umap.find(reminder);  
11             if (itr != umap.end()) {  
12                 res.push_back(i);  
13                 res.push_back(itr->second);  
14                 break;  
15             } else {  
16                 umap[nums[i]] = i;  
17             }  
18         }  
19         return res;  
20     }  
21 };
```

# LC Array#2: 121. Best Time to Buy and Sell Stock (I), only one transaction

## 122: ((II) allow to buy and sell same day, no limit of transactions

```
1 class Solution {
2 public:
3     int maxProfit(vector<int>& prices) {
4         int cost = prices[0];
5         int profit = 0;
6         for (int i = 1; i < prices.size(); i++) {
7             if (prices[i] >= cost) {
8                 profit = max(profit, prices[i] - cost);
9             } else {
10                 cost = prices[i]; // starting over if we see lower co
11             }
12         }
13         return profit;
14     }
15 };
```

```
1 class Solution {
2 public:
3     int maxProfit(vector<int>& prices) {
4         int profit = 0;
5         for (int i = 1; i < prices.size(); i++) {
6             if (prices[i] > prices[i-1]) { // for any day, if we seen today price is higher
7                 profit += prices[i] - prices[i-1]; // we always say we bought it yesterday
8             } // if tomorrow is even higher, we said we bought it back yesterday
9         }
10         return profit;
11     }
12 };
```

- a

# LC Array#3: 217. Contains Duplicate (easy) 219: Contains Duplicate II

```
1 class Solution {  
2 public:  
3     bool containsDuplicate(vector<int>& nums) {  
4         unordered_set<int> seen;  
5         for (int i = 0; i < nums.size(); i++) {  
6             if (seen.find(nums[i]) == seen.end()) {  
7                 seen.insert(nums[i]);  
8             } else {  
9                 return true;  
10            }  
11        }  
12        return false;  
13    }  
14};
```

## Example 1:

Input: nums = [1,2,3,1]

Output: true

- Scan thru list, use hashSet to record if not seen

## Example 2:

Input: nums = [1,2,3,4]

Output: false

```
1 class Solution {  
2 public:  
3     bool containsNearbyDuplicate(vector<int>& nums, int k) {  
4         unordered_map<int, int> map;  
5         for (int i = 0; i < nums.size(); i++) {  
6             if (map.find(nums[i]) == map.end()) {  
7                 map[nums[i]] = i;  
8             } else {  
9                 int j = map[nums[i]];  
10                if (i - j <= k)  
11                    return true;  
12                else  
13                    map[nums[i]] = i;  
14            }  
15        }  
16        return false;  
17    }  
18};
```

# LC Array#4: 238. Product of Array Except Self

- One pass starting from beginning to make LEFT product
- Another pass starting from end to make RIGHT product

```
1 class Solution {  
2 public:  
3     vector<int> productExceptSelf(vector<int>& nums) {  
4         vector<int> left(nums.size()), right(nums.size());  
5         int tmp = 1;  
6         left[0] = tmp;  
7         for (int i=1; i < nums.size(); i++){  
8             tmp *= nums[i-1];  
9             left[i] = tmp;  
10        }  
11        tmp = 1;  
12        right[nums.size()-1] = tmp;  
13        for (int i = nums.size()-2; i >= 0; i--){  
14            tmp *= nums[i+1];  
15            right[i] = tmp;  
16        }  
17  
18        vector<int> res(nums.size());  
19        for (int i= 0; i < nums.size(); i++) {  
20            res[i] = left[i]*right[i];  
21        }  
22        return res;  
23    }  
24 };
```



# LC Array#5: 53. Maximum Subarray

```
1  class Solution {  
2  public:  
3      int maxSubArray(vector<int>& nums) {  
4          int maxSum = INT_MIN;  
5          int curSum = 0;  
6          for (int i = 0; i < nums.size(); i++) {  
7              curSum += nums[i];  
8              maxSum = max(maxSum, curSum);  
9              if (curSum < 0) curSum = 0;  
10         }  
11  
12         return maxSum;  
13     }  
14 };
```

- Use two variables: curSum, maxSum
- Reset to curSum back to zero if it is negative

# LC Array#6: 152. Maximum Product Subarray

```
1  class Solution {  
2  public:  
3      int maxProduct(vector<int>& nums) {  
4          int maxRes = INT_MIN;  
5  
6          int product = 1;  
7          for (int i = 0; i < nums.size(); i++) { //left to right  
8              product *= nums[i];  
9              maxRes = max(product, maxRes);  
10             if (product == 0) product = 1;  
11         }  
12         product = 1;  
13         for (int i = nums.size()-1; i >= 0; i--) { // right to left  
14             product *= nums[i];  
15             maxRes = max(product, maxRes);  
16             if (product == 0) product = 1;  
17         }  
18  
19         return maxRes;  
20     }  
21 };
```

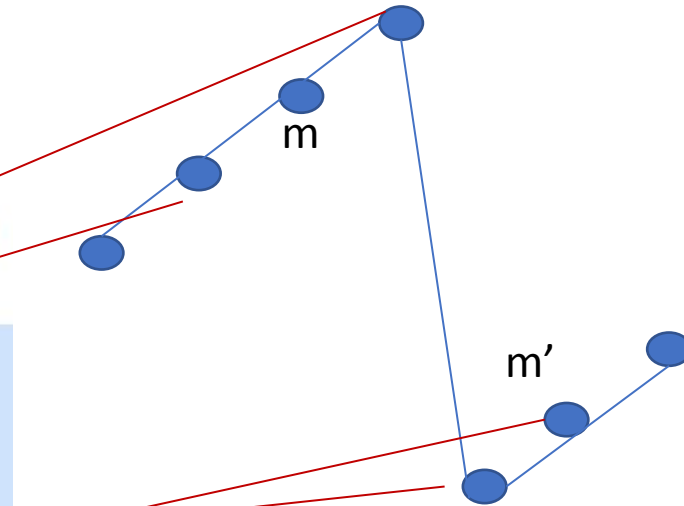
- Use two variables: maxRes, product
- Do product from left to right, track maxRes, if product becomes 0, reset to 1
- Do product from right to left, track maxRes, if product becomes 0, reset to 1

# Binary search: three forms

int

# LC Array#7: 153. Find Minimum in Rotated Sorted Array in $O(\log n)$ time.

```
1 class Solution {
2 public:
3     int findMin(vector<int>& nums) {
4         if (nums.size() == 1) return nums[0];
5         if (nums[0] < nums[nums.size()-1]) return nums[0];
6
7         int l = 0, r = nums.size() - 1;
8         while (l <= r) {
9             if (r == l) return nums[l]; // special handling
10            if (r-l == 1) return (min(nums[l], nums[r])); // special handling
11
12            int m = l + (r-l)/2;
13            if (nums[m] > nums[0]) {
14                if (m+1 < nums.size() && nums[m] > nums[m+1]) {
15                    return nums[m+1];
16                } else {
17                    l = m + 1;
18                }
19            } else {
20                if (m-1 >= 0 && nums[m-1] > nums[m]) {
21                    return nums[m];
22                } else {
23                    r = m - 1;
24                }
25            }
26        }
27        return INT_MIN;
28    }
29};
```



- Special handling when  $r=l$  and  $r-l == 1$  inside loop
- Special handling when  $nums[0]$  is minimum

# LC Array#7.1: 162. Find Peak Element in $O(\log n)$ time.

```
1 class Solution {  
2 public:  
3     int findPeakElement(vector<int>& nums) {  
4         int l = 0, r = nums.size() - 1;  
5         if (r == 0) return 0;  
6         while (l < r) {  
7             if (nums[l] > nums[r]) l = r;  
8             int m = (l+r)/2;  
9             if (nums[m] < nums[m+1]) // climbing  
10                l = m+1;  
11             else  
12                r = m;  
13         }  
14         return l;  
15     }  
16 };
```

**Input:** nums = [1,2,1,3,5,6,4]

**Output:** 5

**Explanation:** Your function can return either index number 1 where the peak element is 2, or index number 5 where the peak element is 6.

A peak element is an element that is strictly greater than its neighbors.

Given an integer array `nums`, find a peak element, and return its index. If the array contains multiple peaks, return the index to **any of the peaks**.

You may imagine that `nums[-1] = nums[n] = -∞`.

You must write an algorithm that runs in  $O(\log n)$  time.



# LC Array#8: 33. Search in Rotated Sorted Array using $O(\log N)$

```
1  class Solution {
2  public:
3      int bst(vector<int>& nums, int l, int r, int target) {
4          int res = -1;
5          while (l <= r) {
6              int m = (l+r)/2;
7              if (nums[m] == target) {
8                  res = m; break;
9              } else {
10                 (nums[m] < target) ? (l=m+1):(r=m-1);
11             }
12         }
13         return res;
14     }
15     int search(vector<int>& nums, int target) {
16         int pivot, l = 0, r = nums.size()-1;
17         if (nums.size() == 1)
18             return nums[0] == target? 0:-1;
19         if (nums[l] < nums[r]) return bst(nums, l, r, target);
20         while (l<=r) {
21             if (l == r) {
22                 pivot = l; break;
23             }
24             if (r - l == 1) {
25                 pivot = nums[l] < nums[r]? l:r;
26                 break;
27             }
28             int m = (l+r)/2;
29             if (m> 0 && (m+1) < (nums.size()) && nums[m+1] > nums[m] && nums[m-1] > nums[m]) {
30                 pivot = m;
31                 break;
32             }
33             if (m>0 && (m+1) < (nums.size()) && nums[m-1] < nums[m] && nums[m] > nums[m+1]) {
34                 pivot = m+1;
35                 break;
36             }
37             if (nums[m] > nums[0])
38                 l = m+1;
39             else
40                 r = m -1;
41         }
42         if (target >= nums[0])
43             return bst(nums, 0, pivot-1, target);
44         else
45             return bst(nums, pivot, nums.size()-1, target);
46     }
47 }
```

- If the array size is 1, return result
- If array is fully sorted, do binary search
- Otherwise find pivot, then decide which segment do perform binary search.

# LC Array#9: 15. 3Sum.

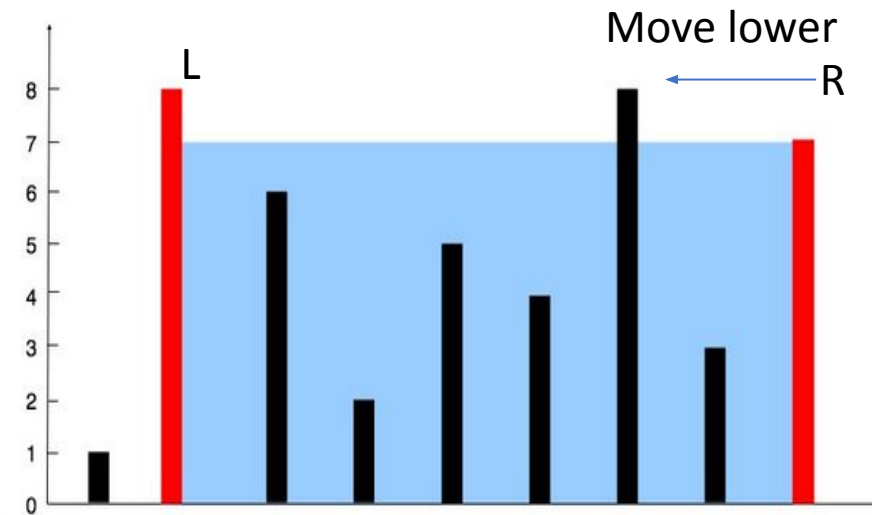
```
1 class Solution {  
2 public:  
3     vector<vector<int>> threeSum(vector<int>& nums) {  
4         vector<vector<int>> res;  
5         if (nums.size() < 3) return res;  
6         sort(nums.begin(), nums.end());  
7         for (int i = 0; i < nums.size()-2; i++) {  
8             if (nums[i] > 0) continue;  
9             if (i > 0 && (nums[i] == nums[i-1])) continue; // avoid dup  
10            int target = -nums[i];  
11            int l = i+1; int r = nums.size()-1;  
12            while (l < r) {  
13                if (nums[l] + nums[r] == target) {  
14                    vector<int> ans;  
15                    ans.push_back(nums[i]);  
16                    ans.push_back(nums[l]);  
17                    ans.push_back(nums[r]);  
18                    res.push_back(ans);  
19                    // don't break, maybe more answer  
20                }  
21                // advance l or r  
22                if (nums[l] + nums[r] <= target) {  
23                    int ll = l;  
24                    while ((l < nums.size() - 1) && (nums[ll] == nums[l])) l++;  
25                } else {  
26                    int rr = r;  
27                    while ((r > 0) && (nums[rr] == nums[r])) r--;  
28                }  
29            }  
30        }  
31        return res;  
32    }  
33 };
```

- Sort array
- Go thru for i, target is  $-\text{nums}[i]$
- Using two pointers between  $i+1$  to  $\text{size}()-1$  to find other two numbers
- $O(n^2)$

# LC Array#10: 11. Container With Most Water.

```
1 class Solution {  
2 public:  
3     int maxArea(vector<int>& height) {  
4         int res = INT_MIN;  
5  
6         #ifdef BRUTEFORCE  
7         for (int i = 0; i < height.size() - 1; i++) {  
8             for (int j = i+1; j < height.size(); j++) {  
9                 int area = (j-i)*min(height[i], height[j]);  
10                res = max(res, area);  
11            }  
12        }  
13        #endif  
14        // two pointers  
15        int l = 0, r = height.size() - 1;  
16        while (l < r) {  
17            int a = (r-l)*min(height[l], height[r]);  
18            res = max(res, a);  
19            if (height[l] < height[r]) {  
20                l++;  
21            } else {  
22                r--;  
23            }  
24        }  
25        return res;  
26    }  
27 };
```

Example 1:



Input: height = [1,8,6,2,5,4,8,3,7]

Output: 49

Explanation: The above vertical lines are represented by array [1,8,6,2,5,4,8,3,7]. In this case, the max area of water (blue section) the container can contain is 49.

# LC#528 PrefixSum

## 528. Random Pick with Weight

Medium 1494 3131 Add to List Share

You are given an array of positive integers  $w$  where  $w[i]$  describes the weight of  $i^{\text{th}}$  index (0-indexed).

We need to call the function `pickIndex()` which **randomly** returns an integer in the range  $[0, w.length - 1]$ . `pickIndex()` should return the integer proportional to its weight in the  $w$  array. For example, for  $w = [1, 3]$ , the probability of picking the index 0 is  $1 / (1 + 3) = 0.25$  (i.e 25%) while the probability of picking the index 1 is  $3 / (1 + 3) = 0.75$  (i.e 75%).

More formally, the probability of picking index  $i$  is  $w[i] / \text{sum}(w)$ .

```
1 class Solution {
2     vector<int> prefixSum;
3     int N;
4 public:
5     Solution(vector<int>& w) {
6         N = w.size();
7         prefixSum.push_back(w[0]);
8         for (int i = 1; i < N; i++) {
9             prefixSum.push_back(prefixSum.back()+w[i]);
10        }
11    }
12
13    int pickIndex() {
14        float r = ((float)rand()) / RAND_MAX; // get a number between 0-1
15        int target = (int) (r * prefixSum[N-1]);
16        for (int i = 0; i < N; i++) {
17            if (target < prefixSum[i]) return i;
18        }
19        return N-1;
20    }
21 };|
22
```

# Bit Manipulation

## Types of Bitwise Operators

Operator	Name	Example	Result
&	Bitwise AND	6 & 3	2
	Bitwise OR	10   10	10
^	Bitwise XOR	2 ^ 2	0
~	Bitwise 1's complement	~9	-10
<<	Left-Shift	10 << 2	40
>>	Right-Shift	10 >> 2	2



# LC BITS#1: 371. Sum of Two Integers.

```
1 class Solution {  
2 public:  
3     int getSum(int a, int b) {  
4         long mask = 0xFFFFFFFF;  
5         while (b) {  
6             int sum = a ^ b;  
7             int carry = ((a & b) & mask) << 1; // avoid negative# shift error  
8             a = sum;  
9             b = carry;  
10        }  
11        return a;  
12    }  
13 };  
14  
15
```

- $(A \& B) \ll 1 \Rightarrow$  carry,  $A \wedge B \Rightarrow$  answer
- Use **long** mask=0xFFFFFFFF

# LC BITS#2: 268. Missing Number.

```
1 class Solution {  
2 public:  
3     int missingNumber(vector<int>& nums) {  
4         int res = nums.size();  
5         for (int i = 0; i < nums.size(); i++) {  
6             res ^= i;  
7             res ^= nums[i];  
8         }  
9         return res;  
10    }  
11 };
```

Given an array `nums` containing `n` distinct numbers in the range `[0, n]`, return *the only number in the range that is missing from the array*.

**Follow up:** Could you implement a solution using only  $O(1)$  extra space complexity and  $O(n)$  runtime complexity?

## Example 1:

Input: `nums = [3,0,1]`

Output: 2

Explanation: `n = 3` since there are 3 numbers, so all numbers are in the range `[0,3]`. 2 is the missing number in the range since it does not appear in `nums`.

- $(A \& B) \ll 1 \Rightarrow \text{carry}$ ,  $A \wedge B \Rightarrow \text{answer}$

# LC BITS#3: 190. Reverse Bits.

```
1 class Solution {  
2 public:  
3     uint32_t reverseBits(uint32_t n) {  
4         unsigned int l_mask = 1 << 31, r_mask = 1;  
5         for (int i = 0; i < 16; i++) {  
6             int r_ = (n & r_mask);  
7             int l_ = (n & l_mask);  
8             n = (n & ~l_mask) | (r_?l_mask:0);  
9             n = (n & ~r_mask) | (l_?r_mask:0);  
10            l_mask = l_mask >>1 ;  
11            r_mask = r_mask << 1;  
12        }  
13        return n;  
14    }  
15 };
```

- a

# LC BITS#4: counting 1 bit of an interger.

```
while (n) { counter++; n &=(n-1);}
__builtin_popcount(n);
```

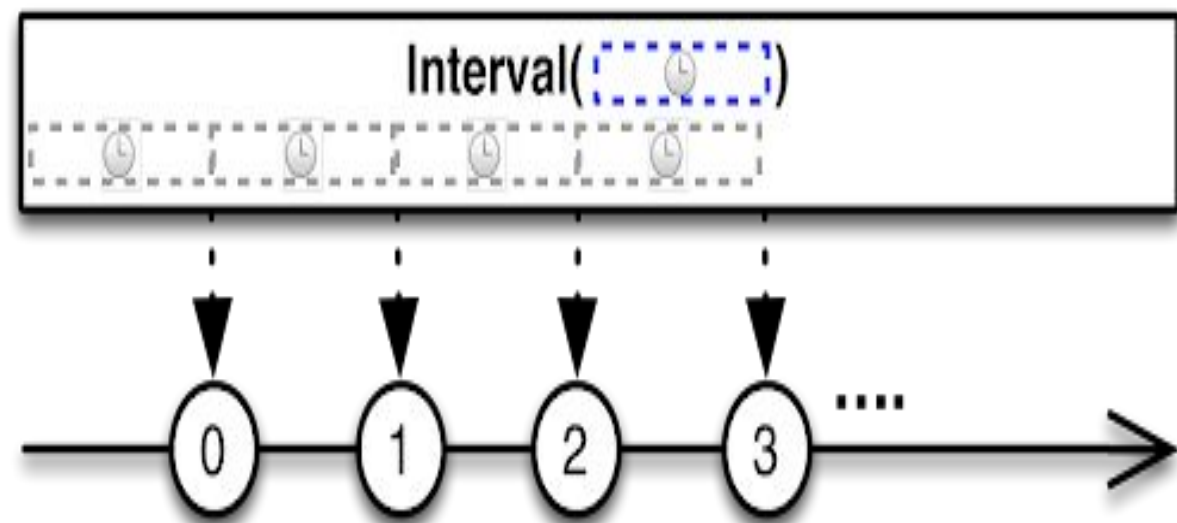
- `__builtin_popcount()`: count number of 1 bit, leverage “**popcnt**” instruction
- `__builtin_clz()`: count leading zero, “lea” instruction
- `__builtin_ctz()`: count trailing zero, “**tzcnt**” instruction
- `__builtin_parity()`: parity check

# LC BITS#5: check if a number is power of 4.

```
bool isPowerOfFour(int n) {  
    return !(n&(n-1)) && (n&0x55555555);  
    //check the 1-bit location;  
}
```

- Check it has only one bit and bit location is  
0x0101, 0101, 0101, 0101,0101,0101,0101,0101
- If it is power of 8, which is  $(2*2*2)^N$ , bit  
location:  
0x0100,1001,0010,0100,1001,0010,0100,1001  
0x49249249





# LC INTERVAL#1: 57. Insert Interval.

```
1  class Solution {
2  public:
3      vector<vector<int>> insert(vector<vector<int>>& intervals, vector<int>& newInterval) {
4          //insert newInterval into intervals to keep ascending order
5          bool inserted = false;
6          for (int i = 0; i < intervals.size(); i++) {
7              if (newInterval[0] <= intervals[i][0]) {
8                  intervals.insert(intervals.begin()+i, newInterval);
9                  inserted = true;
10                 break;
11             }
12         }
13         if (!inserted) intervals.push_back(newInterval);
14
15         stack<vector<int>> st;
16         st.push(intervals[0]);
17         for (int i = 1; i < intervals.size(); i++) {
18             vector<int> top = st.top();
19             if (intervals[i][0] > top[1]) { // no-overlap,
20                 st.push(intervals[i]);
21             } else {
22                 if (intervals[i][1] > top[1]) { // bigger than current interval
23                     top[1] = intervals[i][1]; // extend current interval
24                     st.pop();
25                     st.push(top);
26                 }
27             }
28         }
29         vector<vector<int>> res;
30         while (!st.empty()) {
31             res.insert(res.begin(), st.top());
32             st.pop();
33         }
34         return res;
35     }
36 };
```

Example 1:

Input: intervals = [[1,3],[6,9]], newInterval = [2,5]

Output: [[1,5],[6,9]]

- Given interval list is in ascending order
- Insert newInterval into given list in ascending order
- Merge all intervals in the list using **STACK**

# LC INTERVAL#2: 56. Merge Intervals

Example 1:

Input: intervals = [[1,3],[2,6],[8,10],[15,18]]

Output: [[1,6],[8,10],[15,18]]

Explanation: Since intervals [1,3] and [2,6] overlaps, merge them into [1,6].

```
1 class Solution {
2 public:
3     static bool cmp(vector<int> v1, vector<int> v2) {
4         return v1[0] < v2[0];
5     }
6     vector<vector<int>> merge(vector<vector<int>>& intervals) {
7         if (intervals.size() == 0) return res;
8         sort(intervals.begin(), intervals.end(), cmp);
9
10        vector<vector<int>> res;
11        stack<vector<int>> st;
12        st.push(intervals[0]);
13        for (int i = 1; i < intervals.size(); i++) {
14            vector<int> top = st.top();
15            int c_s = top[0];
16            int c_e = top[1];
17            int n_s = intervals[i][0];
18            int n_e = intervals[i][1];
19            if (n_s <= c_e) {
20                // merge
21                top[1] = n_e > c_e ? n_e : c_e;
22                st.pop();
23                st.push(top);
24            } else {
25                // no overlap, simply push
26                st.push(intervals[i]);
27            }
28        }
29
30        while (!st.empty()) {
31            res.insert(res.begin(), st.top());
32            st.pop();
33        }
34        return res;
35    }
36};
```

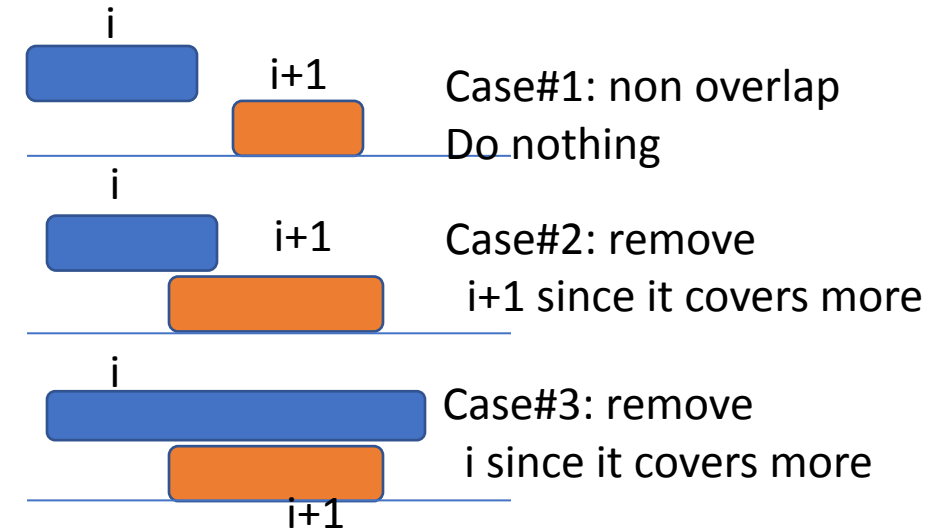
- First sort intervals by using start value in ascending order
- Use stack to merge any overlapped intervals

# LC INTERVAL#3: 435. Non-overlapping Intervals.

Given an array of intervals `intervals` where `intervals[i] = [starti, endi]`, return the minimum number of intervals you need to remove to make the rest of the intervals non-overlapping.

```
1  class Solution {
2  public:
3      static bool cmp(vector<int> v1, vector<int> v2) {
4          return v1[0] < v2[0];
5      }
6      // first sort the intervals based on starting value in ascending order
7      // if interval[i] overlaps with interval[i+1]
8      //     if interval[i+1] end is greater than interval[i], remove i+1
9      //     else remove interval[i], which is longer
10     int eraseOverlapIntervals(vector<vector<int>>& intervals) {
11         if (intervals.size() < 2) return 0;
12
13         sort(intervals.begin(), intervals.end(), cmp);
14         vector<int> prev = intervals[0];
15         int removals = 0;
16         for (int i = 1; i < intervals.size(); i++) {
17             vector<int> cur = intervals[i];
18             if (cur[0] >= prev[1]) {
19                 prev = cur;
20                 continue; // non overlap
21             } else {
22                 if (cur[1] > prev[1]) {
23                     // remove cur
24                     removals++;
25                     // prev pointer stays
26                 } else {
27                     prev = cur;
28                     removals++;
29                 }
30             }
31         }
32         return removals;
33     }
34 }
```

- Sort all interval lists in ascending order
- Go thru interval list, there are three cases below:



# LC INTERVAL#4: 252. Meeting Rooms.

Given an array of meeting time intervals where  $\text{intervals}[i] = [\text{start}_i, \text{end}_i]$ , determine if a person could attend all meetings.

```
1 class Solution {  
2     public:  
3         static bool cmp(vector<int> v1, vector<int> v2) {  
4             return v1[0]<v2[0];  
5         }  
6  
7         bool canAttendMeetings(vector<vector<int>>& intervals) {  
8             if (intervals.size() < 2) return true;  
9  
10            sort(intervals.begin(), intervals.end(), cmp);  
11  
12            vector<int> prev = intervals[0];  
13            for (int i = 1; i < intervals.size(); i++) {  
14                vector<int> cur = intervals[i];  
15                if (prev[1] > cur[0])  
16                    return false;  
17                else {  
18                    prev = cur;  
19                }  
20            }  
21            return true;  
22        }  
23    }  
24};
```

## Example 1:

**Input:** intervals = [[0,30],[5,10],[15,20]]

**Output:** false

- Sort interval by starting time in ascending order
- Check if there is any overlap



# LC INTERVAL#5: 253: Meeting Rooms II.

Example 1:

Input: intervals = [[0,30],[5,10],[15,20]]

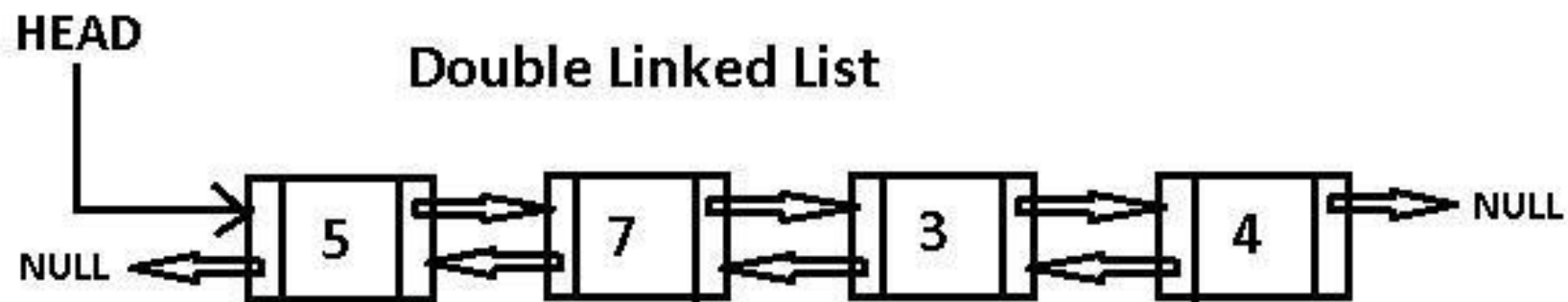
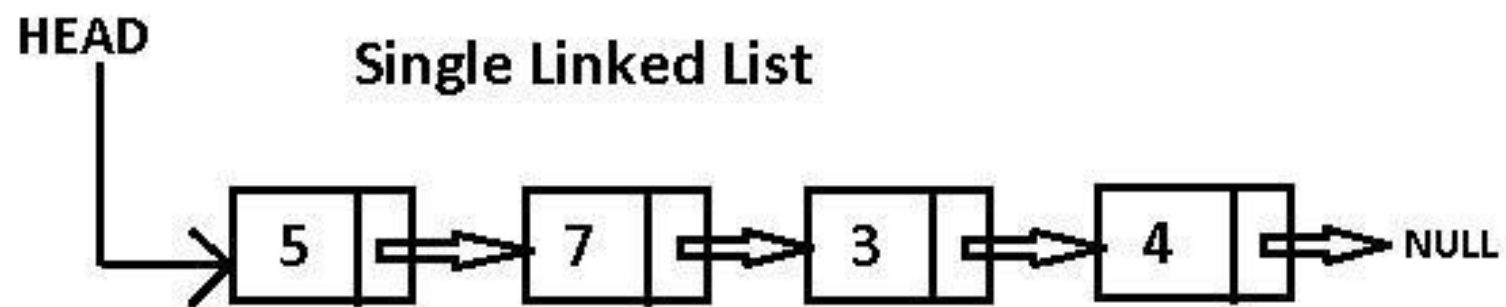
Output: 2

Given an array of meeting time intervals intervals where intervals[i] = [starti, endi], return the minimum number of conference rooms required.

```
1 class Solution {
2 public:
3     static bool cmp(int i1, int i2) {
4         if (abs(i1) == abs(i2)) {
5             return i1 < i2;
6         }
7         return abs(i1) < abs(i2);
8     }
9
10    int minMeetingRooms(vector<vector<int>>& intervals) {
11        vector<int> times;
12        for (int i = 0; i < intervals.size(); i++) {
13            times.push_back(intervals[i][0]);
14            times.push_back(-intervals[i][1]);
15        }
16        sort(times.begin(), times.end(), cmp);
17        int res = 0;
18        int rooms = 0;
19        for (int i = 0; i < times.size(); i++) {
20            if (times[i] >= 0) {
21                rooms++;
22                res = rooms > res ? rooms : res;
23            } else {
24                rooms--;
25            }
26        }
27        return res;
28    }
29};
```

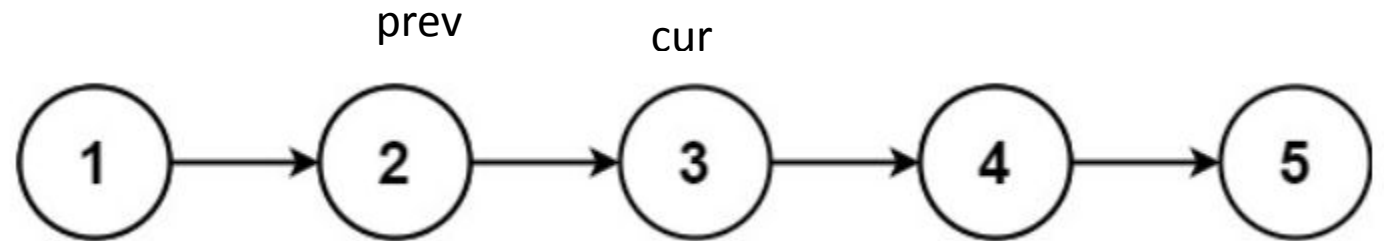
- Merge all starting time and ending time into one single vector, but ending time as negative value
- Sort vector using absolute value, if it is equal, negative first
- Go thru vector list, see positive time, increase 1,
- See negative time, decrease one.
- Record the max rooms





# LC LinkedList#1: 206. Reverse Linked List

```
1 ▾ /**
2   * Definition for singly-linked list.
3   * struct ListNode {
4   *     int val;
5   *     ListNode *next;
6   *     ListNode() : val(0), next(nullptr) {}
7   *     ListNode(int x) : val(x), next(nullptr) {}
8   *     ListNode(int x, ListNode *next) : val(x), next(next) {}
9   * };
10  */
11 ▾ class Solution {
12  public:
13 ▾     ListNode* reverseList(ListNode* head) {
14     ListNode *prev = NULL, *cur = head;
15 ▾     while (cur) {
16         ListNode *tmp = cur->next;
17         cur->next = prev;
18         prev = cur;
19         cur = tmp;
20     }
21     return prev;
22 }
23 };
```



# LC LinkedList#2: 141. Linked List Cycle

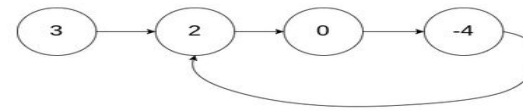
```
9  class Solution {
10  public:
11      bool hasCycle(ListNode *head) {
12          ListNode *s1 = head, *s2;
13          if (!s1) return false;
14          if (s1->next) s2 = s1->next;
15          bool cycle = false;
16          while (s1 && s2) {
17              if (s1 == s2) {
18                  cycle = true;
19                  break;
20              }
21              s1 = s1->next;
22              s2 = s2->next? s2->next->next:NULL;
23          }
24          return cycle;
25      }
26  };
```

Given `head`, the head of a linked list, determine if the linked list has a cycle in it.

There is a cycle in a linked list if there is some node in the list that can be reached again by continuously following the `next` pointer. Internally, `pos` is used to denote the index of the node that tail's `next` pointer is connected to. **Note that `pos` is not passed as a parameter.**

Return `true` if there is a cycle in the linked list. Otherwise, return `false`.

**Example 1:**



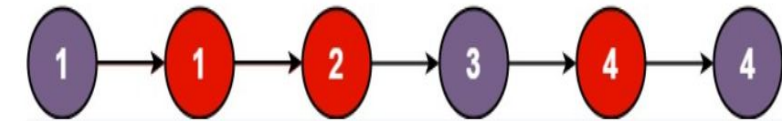
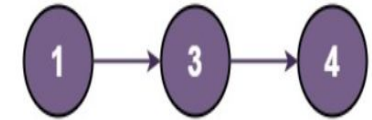
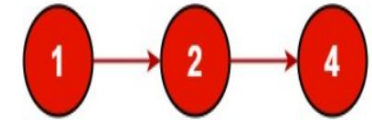
- Use two pointers: s1 move with one step, s2 move two steps
- If there is loop, s1 and s2 will never NULL and will be s1==s2 eventually.
- Or use **hashSet "seen"**

```
9  class Solution {
10  public:
11      ListNode *detectCycle(ListNode *head) {
12          if (head == NULL) return NULL;
13          unordered_set<ListNode *> seen;
14          ListNode *s1 = head;
15          while (s1) {
16              if (seen.find(s1) != seen.end()) return s1;
17              seen.insert(s1);
18              s1 = s1->next;
19          }
20          return NULL;
21      }
22  };
```

# LC LinkedList#3: 21. Merge Two Sorted Lists.

```
11 class Solution {
12 public:
13     ListNode* mergeTwoLists(ListNode* l1, ListNode* l2) {
14         ListNode *head = NULL, *tmp, *cur;
15         if (!l1) return l2;
16         if (!l2) return l1;
17
18         while (l1 && l2) {
19             if (l1->val <= l2->val) {
20                 tmp = l1;
21                 l1 = l1->next;
22             } else {
23                 tmp = l2;
24                 l2 = l2->next;
25             }
26             if (!head) {
27                 head = tmp;
28                 cur = tmp;
29             } else {
30                 cur->next = tmp;
31                 cur = tmp;
32             }
33         }
34         if (l1) {
35             cur->next = l1;
36         }
37         if (l2) {
38             cur->next = l2;
39         }
40
41         return head;
42     }
43 }
```

Example 1:



Input: l1 = [1,2,4], l2 = [1,3,4]

Output: [1,1,2,3,4,4]

- Use “cur” to hold current latest merged node
- Use “head” to hold return value
- Compare l1 and l2, advance the smaller value pointer
- If one becomes NULL, simply connect the other remaining

# LC LinkedList#4: 23. Merge k Sorted Lists.(HARD)

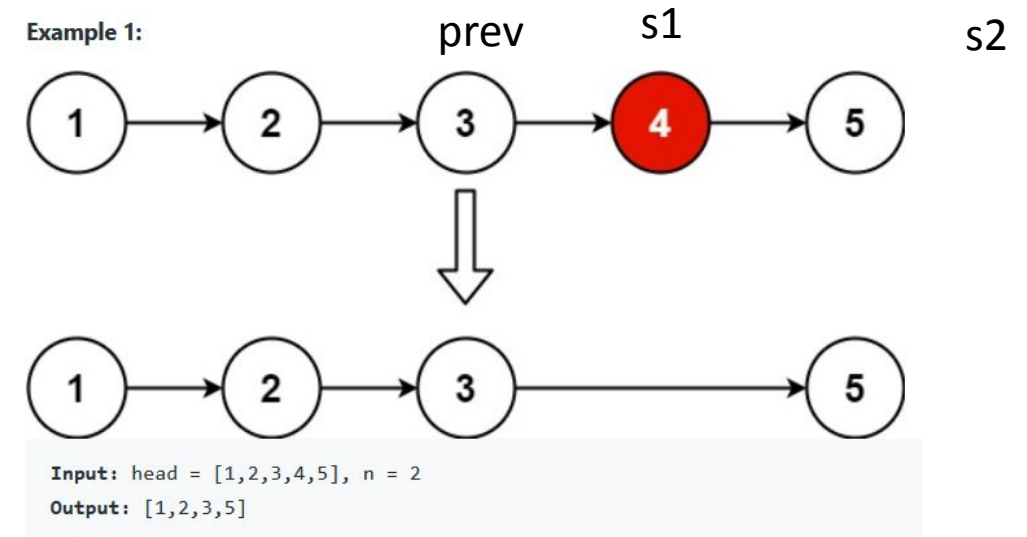
```
11 class Cmp {
12     public:
13     bool operator()(ListNode* n1, ListNode* n2) {
14         return n1->val > n2->val; // ascending order
15     }
16 };
17 class Solution {
18     public:
19     static bool cmp(ListNode* n1, ListNode *n2) {
20         return n1->val < n2->val;
21     }
22
23     ListNode* mergeKLists(vector<ListNode*>& lists) {
24         priority_queue<ListNode *, vector<ListNode *>, Cmp> miniHeap;
25         for (int i = 0; i < lists.size(); i++) {
26             ListNode *head = lists[i];
27             while (head) {
28                 miniHeap.push(head);
29                 head = head->next;
30             }
31         }
32         ListNode *root = NULL, *cur = NULL, *tmp;
33         while (!miniHeap.empty()) {
34             ListNode *tmp = miniHeap.top();
35             tmp->next = NULL;
36             if (!root) {
37                 root = tmp;
38                 cur = root;
39             } else {
40                 cur->next = tmp;
41                 cur = tmp;
42             }
43             miniHeap.pop();
44         }
45         return root;
46     }
47 };|
```

- Create a miniHeap using priority\_queue (or simply use vector sorting)
- Traversal all K lists, add all nodes into minheap
- Pop minheap node one by one and form a new list to return



# LC LinkedList#5: 19. Remove Nth Node From End of List.

```
11 class Solution {
12 public:
13     ListNode* removeNthFromEnd(ListNode* head, int n) {
14         ListNode *s1 = head, *s2 = head, *prev = NULL;
15         for (int i = 0; i < n; i++) {
16             s2 = s2->next;
17         }
18         while (s1 && s2) {
19             prev = s1;
20             s1 = s1->next;
21             s2 = s2->next;
22         }
23         if (prev)
24             prev->next = s1->next;
25         else
26             return s1->next;
27
28         return head;
29     }
30 };
```



- Use two pointers: s2 is N step ahead of s1
- When s2 becomes NULL, S1 is the node to be removed
- Use “prev” hold node before s1

# LC LinkedList#6: 143. Reorder List.

```
11 class Solution {
12 public:
13     void reorderList(ListNode* head) {
14         vector<ListNode*> nodes;
15         ListNode* cur = head;
16         while (cur) { // put all nodes into vector list with next set to NULL
17             ListNode* tmp = cur->next;
18             cur->next = NULL;
19             nodes.push_back(cur);
20             cur = tmp;
21         }
22         ////////////////
23         int N = nodes.size() - 1;
24         cur = NULL;
25         for (int i = 0; i < N; i++) {
26             if (i < (N-i)) {
27                 nodes[i]->next = nodes[N-i];
28                 if (cur)
29                     cur->next = nodes[i];
30                 cur = nodes[N-i];
31             } else if (i == (N-i)) { // point to same node
32                 cur->next = nodes[i]; // last node
33             } else {
34                 break;
35             }
36         }
37     }
38 };
```

You are given the head of a singly linked-list. The list can be represented as:

$$L_0 \rightarrow L_1 \rightarrow \dots \rightarrow L_{n-1} \rightarrow L_n$$

Reorder the list to be on the following form:

$$L_0 \rightarrow L_n \rightarrow L_1 \rightarrow L_{n-1} \rightarrow L_2 \rightarrow L_{n-2} \rightarrow \dots$$

You may not modify the values in the list's nodes. Only nodes themselves may be changed.

- Push all ListNode into vector list
- Take out node from vector as required to form a new list.

**m-by-n** matrix

---

$a_{i,j}$

**n columns**  $\xrightarrow{\text{j changes}}$

**m rows**

$\downarrow \text{i changes}$


$$\begin{bmatrix} a_{1,1} & a_{1,2} & a_{1,3} & \dots \\ a_{2,1} & a_{2,2} & a_{2,3} & \dots \\ a_{3,1} & a_{3,2} & a_{3,3} & \dots \\ \vdots & \vdots & \vdots & \ddots \end{bmatrix}$$

# LC Matrix#1: 73. Set Matrix Zeroes.

```
1 class Solution {
2 public:
3     void setZeroes(vector<vector<int>>& matrix) {
4         int R = matrix.size();
5         int C = matrix[0].size();
6         unordered_set<int> r_zero;
7         unordered_set<int> c_zero;
8
9         for (int r = 0; r < R; r++) {
10             vector<int> row = matrix[r];
11             for (int c = 0; c < C; c++) {
12                 if (matrix[r][c] == 0) {
13                     if (r_zero.find(r) == r_zero.end()) {
14                         // set this row to zero
15                         r_zero.insert(r);
16                     }
17                     if (c_zero.find(c) == c_zero.end()) {
18                         // set this column to zero
19                         c_zero.insert(c);
20                     }
21                 }
22             }
23         }
24         for (int r = 0; r < R; r++) {
25             vector<int> row = matrix[r];
26             for (int c = 0; c < C; c++) {
27                 if (r_zero.find(r) != r_zero.end()
28                     || c_zero.find(c) != c_zero.end()) {
29                     matrix[r][c] = 0;
30                 }
31             }
32         }
33     }
34 };
```

Example 1:

1	1	1
1	0	1
1	1	1



1	0	1
0	0	0
1	0	1

**Input:** matrix = [[1,1,1],[1,0,1],[1,1,1]]

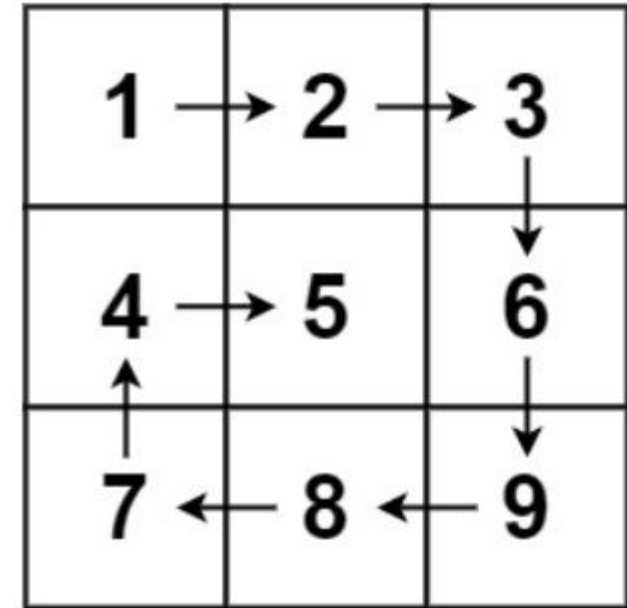
**Output:** [[1,0,1],[0,0,0],[1,0,1]]

- Scan thru matrix to mark down rows and columns (into hashset) to be set to zero
- Second time go thru matrix again and check hashSets , either r or c in the its set, set matrix cell to be zero

# LC Matrix#2: 54. Spiral Matrix.

```
1  class Solution {
2  public:
3      vector<int> spiralOrder(vector<vector<int>>& matrix) {
4          vector<int> res;
5          vector<vector<int>> DIRS {{0, 1}, {1, 0}, {0, -1}, {-1, 0}};
6          int R = matrix.size();
7          int C = matrix[0].size();
8          int r = 0, c = 0, mov_dir = 0;
9          int VISITED = INT_MAX;
10         while (1) {
11             int r_off = DIRS[mov_dir][0];
12             int c_off = DIRS[mov_dir][1];
13             res.push_back(matrix[r][c]);
14             matrix[r][c] = VISITED;
15             int nr = r + r_off;
16             int nc = c + c_off;
17             if (nr < 0 || nr >= R || nc < 0 || nc >= C || matrix[nr][nc] == INT_MAX) {
18                 // time to change direction
19                 mov_dir = (mov_dir + 1) % 4;
20                 r = r + DIRS[mov_dir][0];
21                 c = c + DIRS[mov_dir][1];
22                 if (r < 0 || r >= R || c < 0 || c >= C) break;
23                 if (matrix[r][c] == INT_MAX) break;
24             } else {
25                 r = nr;
26                 c = nc;
27             }
28         }
29         return res;
30     }
31 };
```

Example 1:



**Input:** matrix = [[1,2,3],[4,5,6],[7,8,9]]

**Output:** [1,2,3,6,9,8,7,4,5]

- Set up four directions: L->R, U->D, R->L, D->U with offset
- Starts with [0][0], each visited node filled with INT\_MAX
- When we move node, change direction when either reach out of boundary or VISITED node



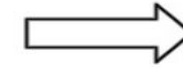
# LC Matrix#3: 48. Rotate Image.

You are given an  $n \times n$  2D matrix representing an image, rotate the image by 90 degrees (clockwise).

```
1  class Solution {
2  public:
3      void rotate(vector<vector<int>>& matrix) {
4          int N = matrix.size();
5          for (int r = 0; r < N; r++) {
6              for (int c = r; c < N; c++) { // column starts with r !!!
7                  int t = matrix[r][c];
8                  matrix[r][c] = matrix[c][r];
9                  matrix[c][r] = t;
10             }
11         }
12
13         // rotate clockwise , reverse each row
14         for (int r = 0; r < N; r++) {
15             int l_ = 0, r_ = N-1;
16             while (l_ < r_) {
17                 int t = matrix[r][l_];
18                 matrix[r][l_] = matrix[r][r_];
19                 matrix[r][r_] = t;
20
21                 l_++; r_--;
22             }
23         }
24         #ifdef ROTATE_COUNTERCLOCK
25         // if rotate counterclockwise, reverse each column
26         for (int c = 0; c < N; c++) {
27             int t_ = 0, b_ = N-1;
28             while (t_ < b_) {
29                 int t = matrix[t_][c];
30                 matrix[t_][c] = matrix[b_][c];
31                 matrix[b_][c] = t;
32
33                 t_++; b_--;
34             }
35         }
36         #endif
37     }
38 }
```

Example 1:

1	2	3
4	5	6
7	8	9



7	4	1
8	5	2
9	6	3

Input: matrix = [[1,2,3],[4,5,6],[7,8,9]]

Output: [[7,4,1],[8,5,2],[9,6,3]]

- Flip value diagnose
- Reverse each row for clockwise
- (Reverse each column for counter-clockwise)



# LC Matrix#4: 79. Word Search.

- Key points:

```
1 class Solution {
2 public:
3     int R = 0;
4     int C = 0;
5     //vector<vector<bool>> map;
6
7     bool backtrack(int r, int c, int pos, string word, vector<vector<char>>& board, vector<vector<bool>> map) {
8         if (pos == word.length()-1) return true;
9
10        if (map[r][c]) return false;
11
12        map[r][c] = true;
13        // check around neighbors (T, D, L, R) if it matches next char, then call recursively
14        vector<vector<int>> DIRS { {-1, 0}, {1, 0}, {0, 1}, {0, -1} };
15        for (int i = 0; i < DIRS.size(); i++) {
16            int nr = DIRS[i][0] + r;
17            int nc = DIRS[i][1] + c;
18            if (nr < 0 || nr >= R || nc < 0 || nc >= C || map[nr][nc]) continue;
19
20            if (board[nr][nc] == word.at(pos+1)) {
21                vector<vector<bool>> nmap = map; // pass new map down
22                if (backtrack(nr, nc, pos+1, word, board, nmap))
23                    return true;
24            }
25        }
26        return false;
27    }
28
29    bool exist(vector<vector<char>>& board, string word) {
30        int r = 0, c = 0, pos = 0;
31        R = board.size();
32        C = board[0].size();
33
34        vector<vector<bool>> map;
35        for (int i = 0; i < R; i++) {
36            vector<bool> t(C, false);
37            map.push_back(t);
38        }
39
40        // find all possible first character to get started
41        for (int r = 0; r < R; r++) {
42            for (int c = 0; c < C; c++) {
43                if (board[r][c] == word.at(0))
44                    if (backtrack(r, c, pos, word, board, map))
45                        return true;
46            }
47        }
48        return false;
49    }
50};
```

Example 1:

A	B	C	E
S	F	C	S
A	D	E	E

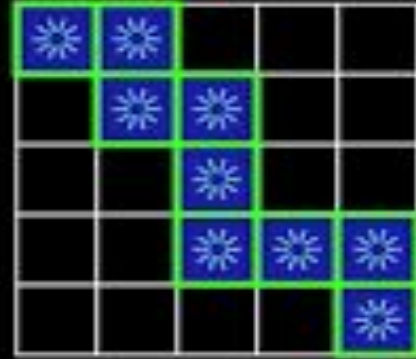
Input: board = [["A","B","C","E"],["S","F","C","S"],["A","D","E","E"]], word = "ABCCED"

Output: true

# LC Matrix#5: TBA.

- $(A \& B) \ll 1 \Rightarrow \text{carry}, A^B \Rightarrow \text{answer}$

# Dynamic Programming



# LC DP#1: 70. Climbing Stairs.

```
1  class Solution {  
2  public:  
3      vector<int> memo;  
4      bool init = true;  
5  int climbStairs(int n) {  
6      if (n <= 2) return n;  
7      if (init) {  
8          memo = vector<int>(n+1, 0);  
9          init = false;  
10     }  
11     if (memo[n] > 0) return memo[n];  
12     int res = climbStairs(n-1) + climbStairs(n-2);  
13     memo[n] = res;  
14     return res;  
15 }
```

- $F(n) = F(n-1) + F(n-2)$  with  $F(1) = 1$ ,  $F(2) = 2$

# LC DP#2: 322. Coin change.

```
3 class Solution {  
4     public:  
5     int coinChange(vector<int>& coins, int amount) {  
6         vector<int> result(amount+1, INT_MAX);  
7         result[0] = 0;  
8         for (int i = 1; i <= amount; i++) {  
9             int ans = INT_MAX-1;  
10            for (int j = 0; j < coins.size(); j++) {  
11                int index = i - coins[j];  
12                if (index >= 0) {  
13                    ans = min(ans, 1 + result[index]);  
14                }  
15            }  
16            result[i] = ans;  
17        }  
18        return result[amount] >= INT_MAX-1? -1: result[amount];  
19    }  
20 }
```

## Example 1:

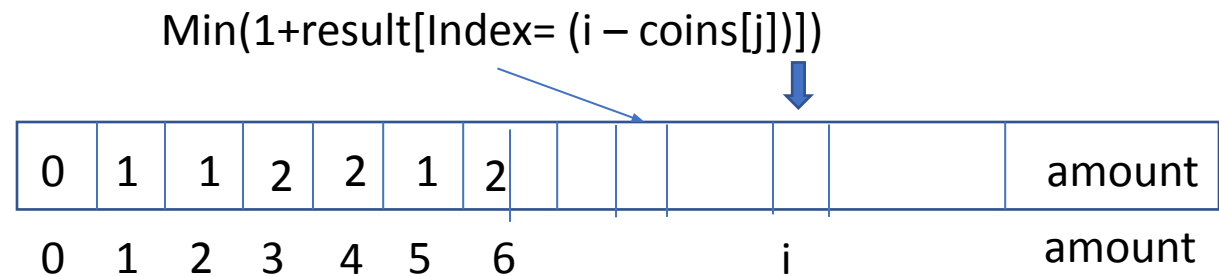
**Input:** coins = [1,2,5], amount = 11

**Output:** 3

**Explanation:** 11 = 5 + 5 + 1

- for each amount, check for each coin
- since coin is always positive, if subtract coin[i] is positive, we use it as index to get a result from that amount, plus 1 (current coin)
- we do for all coins to find the mini number
- Same nature of problem is perfect square LC#279, which is number of [1, 4, 9, 16...]

vector<int> result



# LC DP#3: 300. Longest Increasing Subsequence(LIS) of given array.

Example 1:

**Input:** nums = [10,9,2,5,3,7,101,18]

**Output:** 4

**Explanation:** The longest increasing subsequence is [2,3,7,101], therefore the length is 4.

```
1 class Solution {
2 public:
3     int lengthOfLIS(vector<int>& nums) {
4         vector<int> res(nums.size(), 1);
5
6         for (int i = 1; i < nums.size(); i++) {
7             int ans = 1;
8             for (int j = 0; j <= i; j++) { // going thru all values before me,
9                 if (nums[i] > nums[j]) { // for all values I'm greater
10                     int t = res[j] + 1;
11                     ans = max(ans, t); // keep max number
12                 }
13             }
14             res[i] = ans;
15         }
16         int ret = 0;
17         for (int i = 0; i < nums.size(); i++) ret = max(ret, res[i]);
18         return ret;
19     }
20 };
```

- Start from beginning, for each position, check all values before me
- If it is smaller than me, take its result +1
- Keep max value as result
- If it is Longest Decreasing Subsequence: starting from the end.



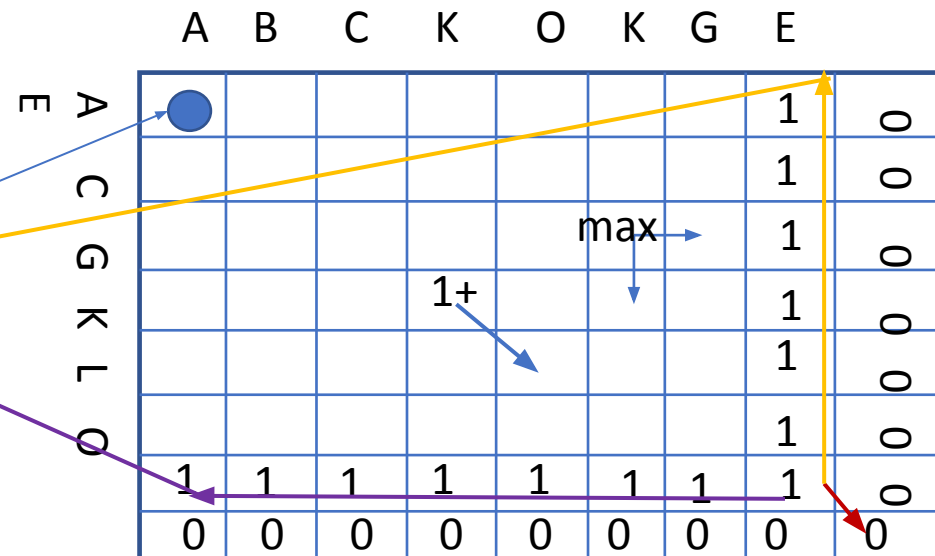
## LC DP#4: 1143. Longest Common Sequence(LCS) of two strings.

```

1 class Solution {
2 public:
3     unordered_map<string, int> dp;
4
5 #ifdef RECURSION
6     int longestCommonSubsequence(string text1, string text2) {
7         if (text1.length() == 0 || text2.length() == 0) return 0;
8         if (dp.find(text1 + "-" + text2) != dp.end()) {
9             return dp[text1 + "-" + text2];
10        }
11        int res = 0;
12
13        string t1 = text1.substr(1);
14        string t2 = text2.substr(1);
15        if (text1.at(0) == text2.at(0)) {
16            res = 1 + longestCommonSubsequence(t1, t2);
17        } else {
18            int lcs1 = longestCommonSubsequence(text1, t2);
19            int lcs2 = longestCommonSubsequence(t1, text2);
20            res = max(lcs1, lcs2);
21        }
22        dp[text1 + "-" + text2] = res;
23        return res;
24    }
25 #endif
26     int longestCommonSubsequence(string text1, string text2) {
27         vector<vector<int>> dp(text1.length()+1, vector<int>(text2.length()+1,0));
28         int r = text1.length()-1;
29         int c = text2.length()-1;
30
31         while (r >= 0 && c >= 0) {
32             for (int c_ = c; c_ >= 0; c_--) {
33                 if (text2.at(c_) == text1.at(r))
34                     dp[r][c_] = 1 + dp[r+1][c_+1];
35                 else
36                     dp[r][c_] = max(dp[r+1][c_], dp[r][c_+1]);
37             }
38             for (int r_ = r; r_ >= 0; r_--) {
39                 if (text2.at(c) == text1.at(r_))
40                     dp[r_][c] = 1 + dp[r_+1][c+1];
41                 else
42                     dp[r_][c] = max(dp[r_+1][c], dp[r_][c+1]);
43             }
44             r--; c--;
45         }
46         return dp[0][0];
47     }
48 };

```

- Use recursion: straightforward, but it takes much more time  $O(2^{\min(N, M)})$
- Basically below plus memo:  
if (text1.at(0) == text2.at(0))  
    return 1 + LCS(text1.substr(1), text2.substr(1));  
else  
    return (max(LCS(text1, text2.substr(1)),  
                LCS(text1.substr(1), text2)));
- Use dp approach: it is  $O(m*n)$



# LC DP#5: 139. word break: given sentence and word dictionary, tell if it can be break using word from dictionary.

```
1 class Solution {
2 public:
3     unordered_map<string, bool> dp;
4
5     bool wordBreak(string s, vector<string>& wordDict) {
6         if (s.length() == 0) return true;
7         if (dp.find(s) != dp.end()) return dp[s];
8
9         // check if any matched word in dict,
10        for (int i = 0; i < wordDict.size(); i++) {
11            string word = wordDict[i];
12            if (s.find(word, 0) == 0) {
13                if (wordBreak(s.substr(word.length()), wordDict)) {
14                    dp[s] = true;
15                    return true;
16                }
17            }
18        }
19        dp[s] = false;
20        return false;
21    }
22};
```

- `std::string::find(substr, pos)`: find first occurrence substr starting index of pos
- `std::string::rfind(substr, pos)`: find last occurrence substr starting from pos
- `Std::string::find_first_of(char c, int pos)`
- Go thru given dictionary, for each word, use `s.find(word, 0)` to see if given string starts with word, if yes, recursive call of substring (after remove word)
- Use memo to improve performance

# Backtracking recipe

```
void Backtrack(res, args)
|
|   if ( GOAL REACHED )
|   |
|   |   add solution to res
|   |
|   |   return
|   |
|   for ( int i = 0; i < NB_CHOICES; i++ )
|   |
|   |   if ( CHOICES[i] is valid )
|   |   |
|   |   |   make choices[i]
|   |   |
|   |   |   Backtrack(res, args)
|   |   |
|   |   |   undo choices[i]
```

# LC DP#6.0 permutation

```
3 // https://www.youtube.com/watch?v=Nabbpl7y4Lo
4 // three steps of backtrack:
5 // #1 check if reaches goal, if yes, save result and return
6 // #2 among all possible choices
7 //   - take valid choice, update in input list
8 //   -- perform backtrack
9 //   - undo the choice, go for next choice
10 //
11 void backtrack(vector<int>& nums, vector<int>& perm, vector<vector<int>>& res, vector<bool>& used) {
12     if (perm.size() == nums.size()) { // check if meets goal
13         res.push_back(perm);         // yes, save result
14         return;
15     }
16     for (int i = 0; i < nums.size(); i++) { // all possible choices
17         if (!used[i]) { // only take valid one
18             used[i] = true; // make the choice
19             perm.push_back(nums[i]); // update parameters
20             backtrack(nums, perm, res, used);
21             perm.pop_back(); // undo the choice
22             used[i] = false; // undo the choice
23         }
24     }
25 }
26
27 vector<vector<int>> permute(vector<int>& nums) {
28     vector<vector<int>> res;
29     vector<int> perm;
30     vector<bool> used(nums.size(), false);
31     backtrack(nums, perm, res, used);
32     return res;
33 }
34 };
```



# LC DP#6: 39. combination sum I

Given an array of **distinct** integers `candidates` and a target integer `target`, return a list of all **unique combinations** of `candidates` where the chosen numbers sum to `target`. You may return the combinations in **any order**.

The **same** number may be chosen from `candidates` an **unlimited number of times**. Two combinations are unique if the frequency of at least one of the chosen numbers is different.

It is **guaranteed** that the number of unique combinations that sum up to `target` is less than 150 combinations for the given input.

```
1 class Solution {
2 public:
3     // reference: https://www.youtube.com/watch?v=yFfv03AE_vA
4     void backtrack(vector<int>& candidates, int start, vector<int> list, vector<vector<int>>& result, int target) {
5         if (target < 0) return;
6         if (target == 0) {
7             if (list.size() > 0) {
8                 vector<int> res1(list);
9                 result.push_back(res1);
10            }
11            return;
12        }
13        for (int i = start; i < candidates.size(); i++) {// start is used to prevent duplicate
14            list.push_back(candidates[i]);
15            backtrack(candidates, i, list, result, target-candidates[i]);
16            list.pop_back();
17        }
18    }
19
20    vector<vector<int>> combinationSum(vector<int>& candidates, int target) {
21        vector<vector<int>> result;
22        vector<int> list;
23
24        backtrack(candidates, 0, list, result, target);
25
26        return result;
27    }
28 };
```

Allow same number to be re-used

- Numbers are distinct
- backtrack allow same number to be re-used

# LC DP#6: 40. combination sum II

Given a collection of candidate numbers ( `candidates` ) and a target number ( `target` ), find all unique combinations in `candidates` where the candidate numbers sum to `target` .

Each number in `candidates` may only be used **once** in the combination.

```
1  class Solution {
2  public:
3      void backtrack(vector<int>& candidates, int start, int target, vector<int> list, vector<vector<int>>& result) {
4          if (target < 0) return;
5          if (target == 0) {
6              vector<int> ans(list);
7              result.push_back(ans);
8              return;
9          }
10         for (int i = start; i < candidates.size(); i++) {
11             // this makes sure number is used only once
12             if (i > start && candidates[i] == candidates[i-1]) continue;
13             if (target-candidates[i] < 0) break; // since it is sorted
14             list.push_back(candidates[i]);
15             backtrack(candidates, i+1, target-candidates[i], list, result); // i+1 to avoid re-use the same number
16             list.pop_back();
17         }
18     }
19
20     vector<vector<int>> combinationSum2(vector<int>& candidates, int target) {
21         vector<vector<int>> result;
22         vector<int> list;
23         sort(candidates.begin(), candidates.end()); // sort in ascending order
24         backtrack(candidates, 0, target, list, result);
25         return result;
26     }
27 };
```

- Values are random, could duplicate
- Sort them in ascending order
- each occurrence of number is used only once
- **Skip the following numbers with same value**



# LC DP#6: 216. combination sum III

Find all valid combinations of  $k$  numbers that sum up to  $n$  such that the following conditions are true:

- Only numbers 1 through 9 are used.
- Each number is used **at most once**.

Return a list of all possible valid combinations. The list must not contain the same combination twice, and the combinations may be returned in any order.

```
1  class Solution {
2  public:
3      void backtrack(int k, int start, int target, vector<int> list, vector<vector<int>> &result) {
4          if (target < 0) return;
5          if (target == 0) {
6              if (list.size() == k) { // result only allow K numbers
7                  vector<int> ans(list);
8                  result.push_back(ans);
9              }
10             return;
11         }
12         if (list.size() >= k) return; // too many stop here
13         for (int i = start; i < 10; i++) { // only use the digits 1-9
14             list.push_back(i);
15             backtrack(k, i+1, target-i, list, result);
16             list.pop_back();
17         }
18     }
19
20     vector<vector<int>> combinationSum3(int k, int n) {
21         vector<vector<int>> result;
22         vector<int> list;
23         backtrack(k, 1, n, list, result);
24         return result;
25     }
26 };
```

- simplified version of II

# LC DP#6: 377. combination sum IV

Given an array of **distinct** integers `nums` and a target integer `target`, return *the number of possible combinations that add up to* `target`.

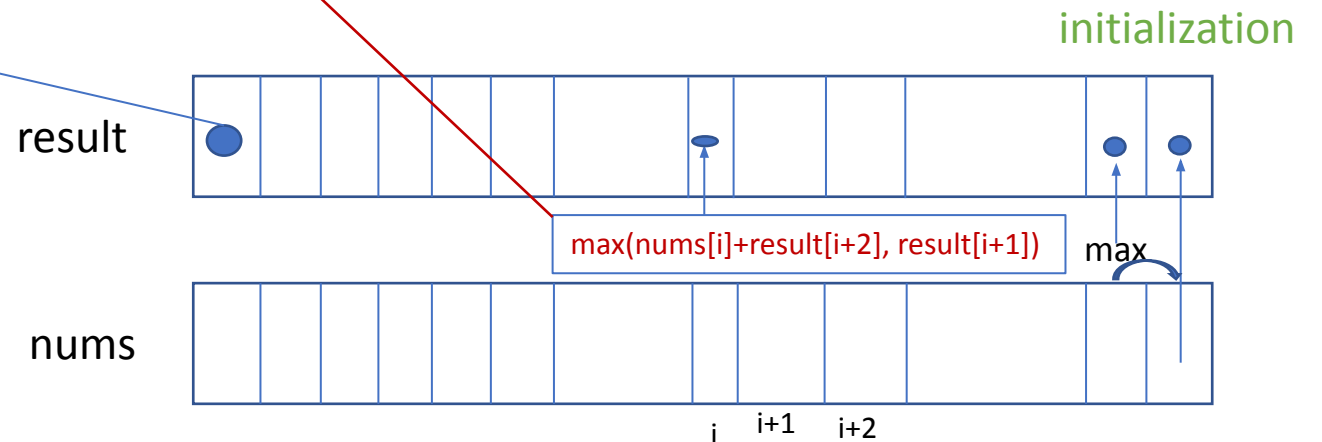
```
1 class Solution {
2     public:
3         unordered_map<string, int> memo;
4         int backtrack(vector<int>& nums, int pos, int target) {
5             int res = 0;
6             if (target < 0) return 0;
7             if (target == 0) {
8                 return 1;
9             }
10            string key = to_string(pos) + "---" + to_string(target);
11            if (memo.find(key) != memo.end()) return memo[key];
12
13            for (int i = pos; i < nums.size(); i++) {
14                if (target - nums[i] < 0) break;
15                res += backtrack(nums, pos, target - nums[i]);
16            }
17            memo[key] = res;
18            return res;
19        }
20
21        int combinationSum4(vector<int>& nums, int target) {
22            int result = 0;
23            sort(nums.begin(), nums.end());
24            result = backtrack(nums, 0, target);
25            return result;
26        }
27    };
--
```

- backtrack with memo of pos+target

# LC DP#7: 198. Rob house: Given an integer array nums representing the amount of money of each house, return the maximum amount of money you can rob tonight without alerting the police.

```
1 class Solution {  
2 public:  
3     int rob(vector<int>& nums) {  
4         if (nums.size() == 1) return nums[0];  
5  
6         vector<int> result(nums.size(), 0);  
7         result[nums.size()-1] = nums[nums.size()-1];  
8         result[nums.size()-2] = max(nums[nums.size()-1], nums[nums.size()-2]);  
9         for (int i = nums.size()-3; i >= 0; i--) {  
10            result[i] = max(result[i+2]+nums[i], result[i+1]);  
11        }  
12        return result[0];  
13    }  
14 }  
15
```

- Since we don't know what is ahead to make it maximum, we start from the end: last one is it self, last second one is if it is greater last, then it is self, otherwise would be last one
- For any given house position "index", compare "its value + next\_next" with "next", take the bigger value



## LC DP#8: 91. decode way. 226 => BBF, BZ, VF

A message containing letters from `A-Z` can be **encoded** into numbers using the following mapping:

```
'A' -> "1"
'B' -> "2"
...
'Z' -> "26"
```

To **decode** an encoded message, all the digits must be grouped then mapped back into letters using the reverse of the mapping above (there may be multiple ways). For example, "11106" can be mapped into:

- "AAJF" with the grouping (1 1 10 6)
- "KJF" with the grouping (11 10 6)

```

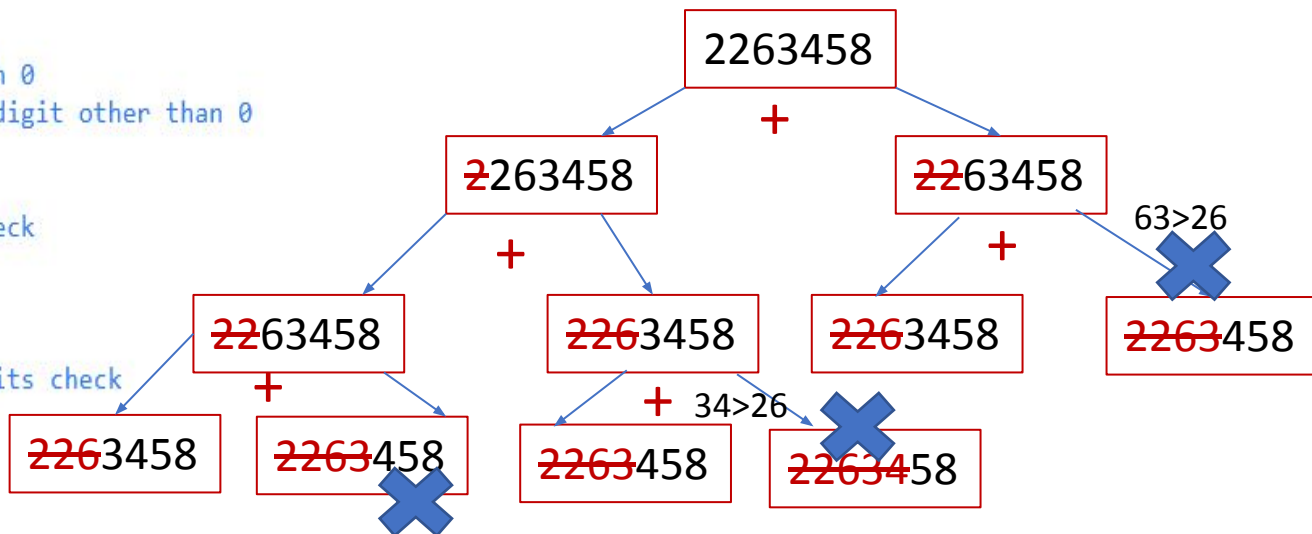
3  class Solution {
4      private :
5      public:
6          unordered_map<string, int> memo;
7
8      int numDecodings(string s) {
9          if (s.length() == 0) return 1;
10         if (s.at(0) == '0') return 0; // can't decode starting with 0
11         if (s.length() == 1) return 1; // good if it is only one digit other than 0
12         if (memo.find(s) != memo.end()) return memo[s];
13
14         int res = numDecodings(s.substr(1)); // spin one digit check
15         if (s.length() >= 2) {
16             int val = stoi(s.substr(0, 2));
17             if (val <= 26) {
18                 res += numDecodings(s.substr(2)); // spin two digits check
19             }
20         }
21         memo[s] = res;
22         return res;
23     }
24 };
25

```

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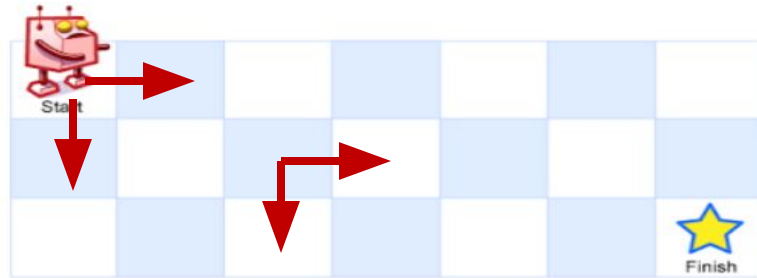
22

- Check special cases:
  - end of string: good (return 1)
  - Starts with 0, can't decode
  - One digit other than 0, good (return 1)
- Check memo cache, if it exists, return result
- Otherwise try to explore with one digit and two digits if possible



# LC DP#9: 62. unique paths/grid walk

Example 1:



Input: m = 3, n = 7

Output: 28

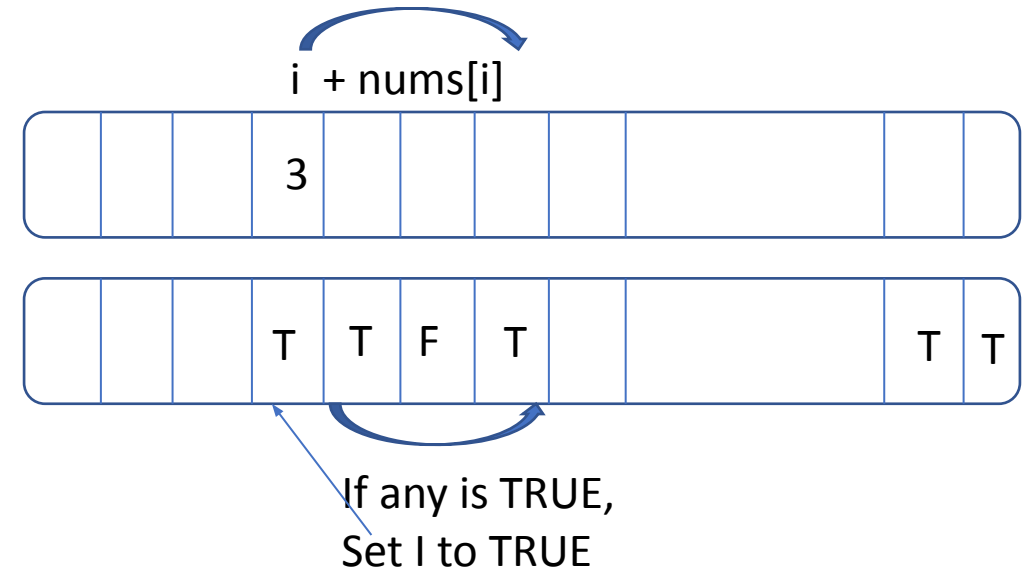
- $Res[m,n] = res[m-1, n] + res[m, n-1]$
- Use a `unordered_map<string, int>` as memo
- Key is "r"=="c"

```
1 class Solution {
2 private:
3     unordered_map<string, int> memo;
4 public:
5
6     int uniquePaths(int m, int n) {
7
8         if (m < 1 || n < 1) return 0;
9         if (m == 1 || n == 1) return 1; // walk bottom or right
10        string key = to_string(m) + "---" + to_string(n);
11        if (memo.find(key) != memo.end()) {
12            return memo[key];
13        }
14        int res = uniquePaths(m-1, n) + uniquePaths(m, n-1);
15        memo[key] = res;
16
17        return res;
18    }
19 };
```



# LC DP#10: 55. Jump Game.

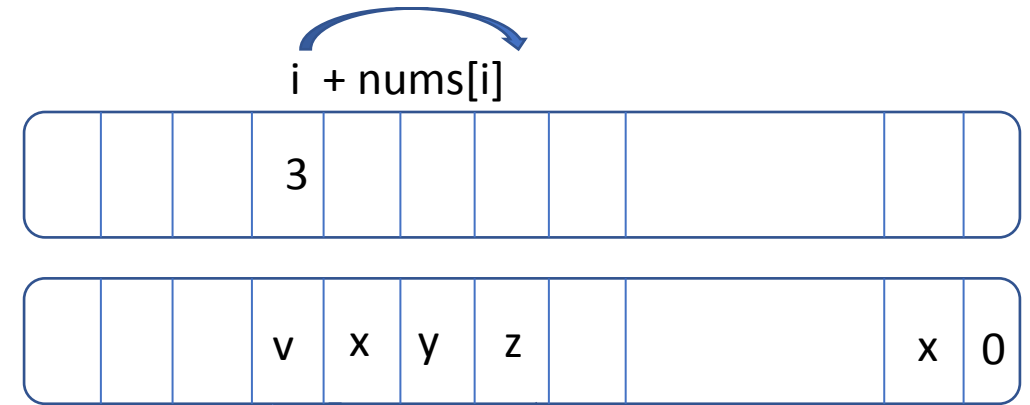
```
class Solution {  
public:  
    bool canJump(vector<int>& nums) {  
        int N = nums.size();  
        vector<bool> status(N, false);  
  
        status[N-1] = true; // last stop is true  
        for (int i = N - 2; i >= 0; i--) {  
            int furestJump = min(i + nums[i], N-1);  
            for (int j = i+1; j <= furestJump; j++) {  
                if (status[j] == true) { // within reach, if any is good  
                    status[i] = true; // we are good.  
                    break;  
                }  
            }  
        }  
        return status[0];  
    }  
};
```





# LC DP#11: 45. Jump Game II.

```
1 class Solution {  
2 public:  
3     int jump(vector<int>& nums) {  
4         int N = nums.size();  
5         vector<int> jumps(N);  
6         jumps[N-1] = 0;  
7  
8         for (int i = N-2; i >= 0; i--) {  
9             int maxPos = min(N-1, i+nums[i]);  
10            int ans = INT_MAX;  
11            for (int j=i+1; j <= maxPos; j++) {  
12                int tmp = jumps[j] != INT_MAX ? (jumps[j] + 1):INT_MAX;  
13                ans = min(ans, tmp);  
14            }  
15            jumps[i] = ans;  
16        }  
17        return jumps[0];  
18    }  
19 };
```



Pick smallest,  
Plus 1, Set to I  
 $V = \min(x, y, \dots, z) + 1$

# LC DP#12: 1306. Jump Game III

Given an array of non-negative integers `arr`, you are initially positioned at `start` index of the array. When you are at index `i`, you can jump to `i + arr[i]` or `i - arr[i]`, check if you can reach to **any** index with value 0.

```
1 class Solution {
2 public:
3     bool helper(vector<int>& arr, int start, vector<bool>& v) {
4         if (arr[start] == 0) return true; // reach desired place
5         if (v[start] == true) return false; // found loop
6
7         v[start] = true; // mark this is visited
8         vector<bool> nv = v;
9         if (start + arr[start] < arr.size()) { // try next round
10             if (helper(arr, start + arr[start], nv)) return true;
11         }
12         if (start - arr[start] >= 0) {
13             if (helper(arr, start - arr[start], nv)) return true;
14         }
15         return false;
16     }
17
18     bool canReach(vector<int>& arr, int start) {
19         vector<bool> v(arr.size(), false);
20
21         return helper(arr, start, v);
22     }
23 }
```

- Recursive call forward and backward if it inbound
- Use a visited map, if reach visited pos, loop detected, NOT possible

## Example 1:

**Input:** `arr = [4,2,3,0,3,1,2]`, `start = 5`

**Output:** `true`

**Explanation:**

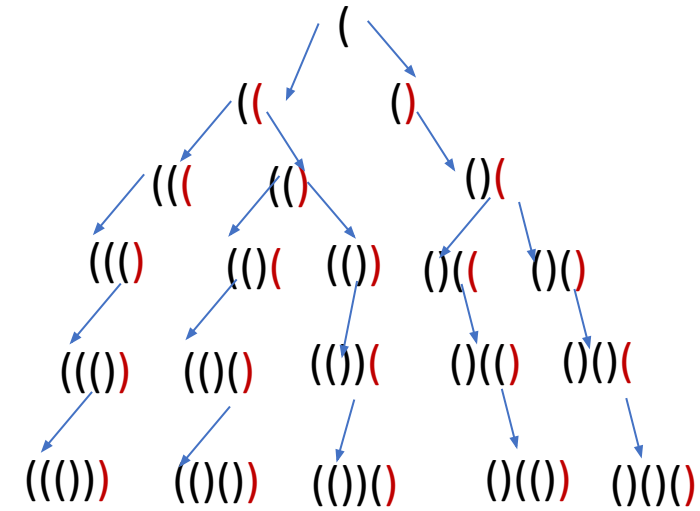
All possible ways to reach at index 3 with value 0 are:

index 5 -> index 4 -> index 1 -> index 3

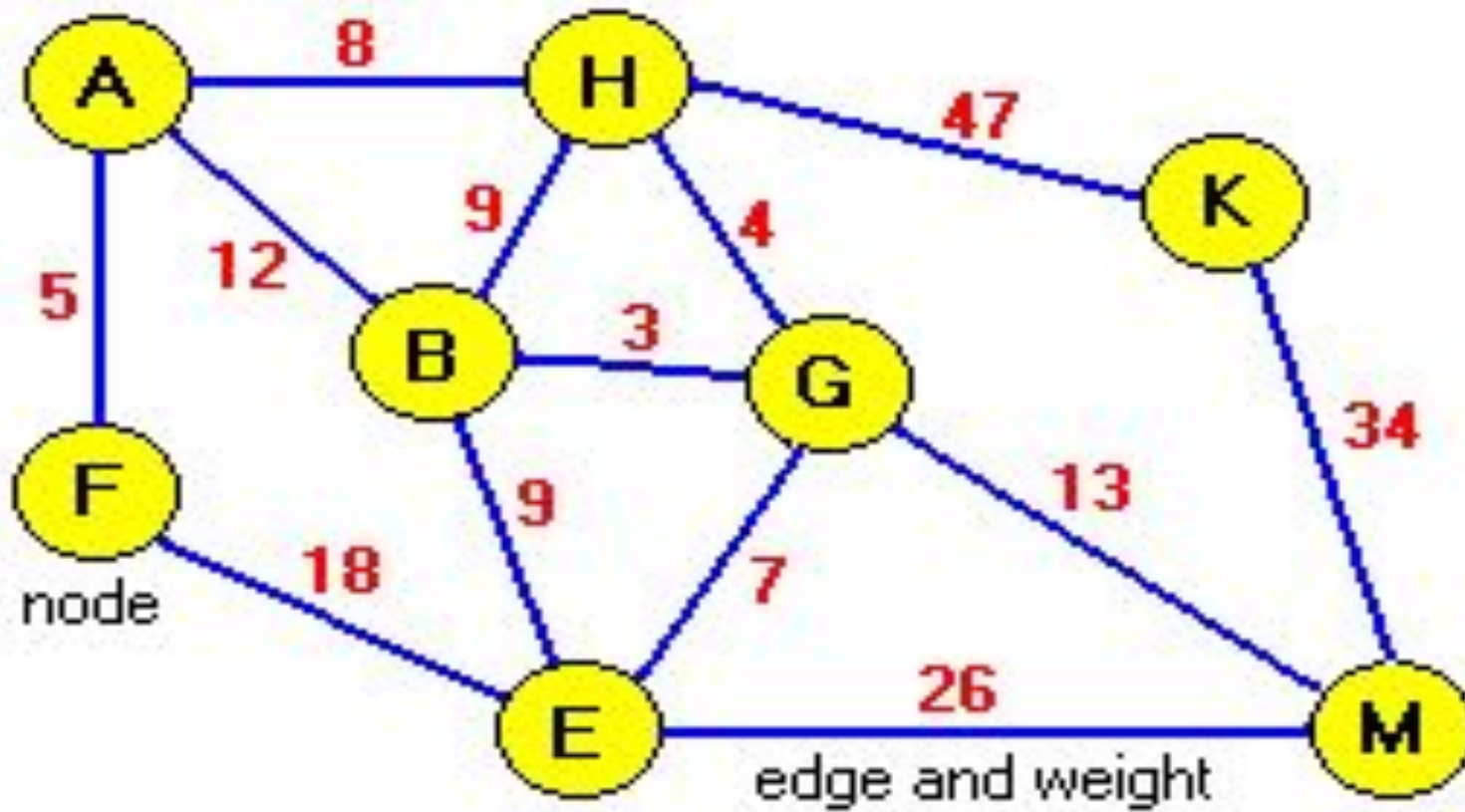
index 5 -> index 6 -> index 4 -> index 1 -> index 3

# LC DP#13: 22: Generate parentheses

```
1 class Solution {  
2 public:  
3     int N;  
4     void backtrack(vector<string>& res, vector<char> cur, int open, int close) {  
5         if (cur.size() == 2*N) {  
6             string ans = "";  
7             for (int i = 0; i < 2*N; i++) ans.append(1,cur[i]);  
8             res.push_back(ans);  
9             return;  
10        }  
11        if (open < N) {  
12            cur.push_back('(');  
13            backtrack(res, cur, open+1, close);  
14            cur.pop_back();  
15        }  
16        if (open > close) {  
17            cur.push_back(')');  
18            backtrack(res, cur, open, close+1);  
19            cur.pop_back();  
20        }  
21    }  
22    vector<string> generateParenthesis(int n) {  
23        vector<string> res;  
24        vector<char> cur;  
25        N = n;  
26        backtrack(res, cur, 0, 0);  
27        return res;  
28    }  
29 };
```



# Graph





# LC Graph#1: 133. clone graph.

```
22 class Solution {
23 public:
24     Node* cloneGraph(Node* node) {
25         Node *root = NULL;
26         queue<Node*> q;
27         unordered_map<int, Node*> newNodes;
28         unordered_set<int> seen;
29         if (node == NULL) return NULL;
30
31         q.push(node);
32
33         while (!q.empty()) {
34             Node *t = q.front();
35             q.pop();
36             Node *n = NULL;
37             if (newNodes.find(t->val) == newNodes.end()) {
38                 n = new Node(t->val);
39                 if (root == NULL) root = n;
40                 newNodes[t->val] = n;
41             } else {
42                 n = newNodes[t->val];
43             }
44
45             for (int i = 0; i < t->neighbors.size(); i++) {
46                 Node *nei = t->neighbors[i];
47                 Node *nnei = NULL;
48                 if (newNodes.find(nei->val) == newNodes.end()) {
49                     nnei = new Node(nei->val);
50                     newNodes[nnei->val] = nnei;
51                 } else {
52                     nnei = newNodes[nei->val];
53                 }
54                 if (find(n->neighbors.begin(), n->neighbors.end(), nnei) == n->neighbors.end())
55                     n->neighbors.push_back(nnei);
56                 if (seen.find(nei->val) == seen.end())
57                     q.push(nei);
58             }
59             seen.insert(n->val);
60         }
61         return root;
62     }
63 };
```

- Use BFS (i.e. use Q) to traversal Graph with help of unordered\_set<int>
- Create a unordered\_map<int, Node\*> to track all created cloned Node, key is node#
- When traverse a node,
  - check cloned node in map, if it doesn't exist, create it.
  - find its all neighbor nodes, check if its cloned node is in map, if not created,
  - Add all cloned neighbors under cloned node

# LC Graph#2: 207. Course Schedule.

There are a total of `numCourses` courses you have to take, labeled from `0` to `numCourses - 1`. You are given an array `prerequisites` where `prerequisites[i] = [ai, bi]` indicates that you **must** take course `bi` first if you want to take course `ai`. Return `true` if you can finish all courses. Otherwise, return `false`.

```
1  class Solution {
2      const int INIT = 0;
3      const int DONE = 1;
4      const int PROCESSING = 2;
5      unordered_map<int, vector<int>> preq_list; // essentially adjacency list of graph
6      vector<int> state;
7  public:
8      bool hasCycle(int course) {
9          if (state[course] == PROCESSING) return true; // cycle detected!
10         state[course] = PROCESSING;
11         if (peq_list.find(course) != preq_list.end()) {
12             vector<int> preqs = preq_list[course];
13             for (int j = 0; j < preqs.size(); j++) {
14                 if (state[preqs[j]] != DONE) {
15                     if (hasCycle(preqs[j]))
16                         return true;
17                 }
18             }
19         }
20         state[course] = DONE;
21         return false;
22     }
23     bool canFinish(int numCourses, vector<vector<int>>& prerequisites) {
24         // go thru to establish preq for each course (node)
25         for (int i = 0; i < prerequisites.size(); i++) {
26             vector<int> elem = prerequisites[i];
27             int course = elem[0];
28             int pre = elem[1];
29             if (peq_list.find(course) != preq_list.end()) {
30                 vector<int> courses = preq_list[course];
31                 courses.push_back(pre);
32                 preq_list[course] = courses; // update
33             } else {
34                 vector<int> preq { pre }; //preq.push_back(pre);
35                 preq_list[course] = preq;
36             }
37         }
38         // now go thru each course, to make sure there is no loop in the preq_list
39         state = vector<int>(numCourses, INIT); // essentially it creates an array, init to zero
40         for (int i = 0; i < numCourses; i++) {
41             if (state[i] == INIT) { // if this course was never initiated, start detect
42                 if (hasCycle(i))
43                     return false;
44             }
45         }
46         return true;
47     }
48 };
```

- Go thru given pre-requisites to establish `peq_list` `unordered_map<int, vector<int>>`: key is course#, value is its pre\_requisite `vector<int>`
- Init state for each course to INIT
- Go thru each course, using GRAPH COLOR to detect if there is any cycle.



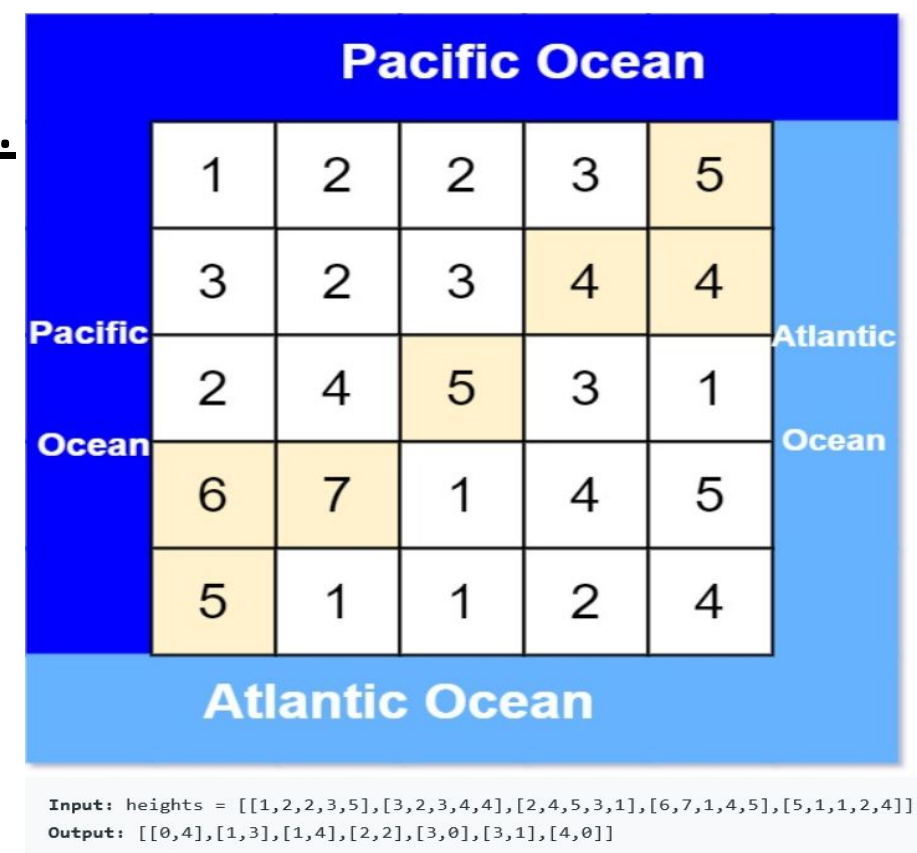
# LC Graph#3: 210. Course Schedule II (topological sort).

```
1 class Solution {
2 public:
3     vector<int> findOrder(int numCourses, vector<vector<int>>& prerequisites) {
4         vector<int> res;
5         vector<vector<int>> preq_list(numCourses);
6         // establish indegree for each course: number of pre-requisites
7         vector<int> indegrees(numCourses, 0);
8         for (int i = 0; i < prerequisites.size(); i++) {
9             int target = prerequisites[i][0];
10            int condition = prerequisites[i][1];
11            preq_list[condition].push_back(target); // collect who requires/needs me
12            indegrees[target] +=1; // indegree table for each course
13        }
14
15        // go thru indegrees table to find initial courses with zero indegree
16        queue<int> q;
17        unordered_set<int> visited;
18        for (int i = 0; i < numCourses; i++) {
19            if (indegrees[i] == 0) {
20                q.push(i);
21                visited.insert(i);
22            }
23        }
24        while (!q.empty()) {
25            int c = q.front(); q.pop();
26            res.push_back(c);
27
28            // find who needs me, decrease indegrees by 1 ,
29            vector<int> targets = preq_list[c];
30            for (int j = 0; j < targets.size(); j++) {
31                indegrees[targets[j]] -=1;
32            }
33            // and add the member with zero into q if it is not seen before
34            for (int i = 0; i < numCourses; i++) {
35                if (indegrees[i] == 0 && visited.find(i) == visited.end()) {
36                    q.push(i);
37                    visited.insert(i);
38                }
39            }
40        }
41        if (res.size() != numCourses) res.clear();
42        return res;
43    }
44};
```

- Go thru given pre-requisites list to establish indegree[target] table and preq\_list<cond, targets> (for a given course, all other courses require/need it)
- Add all course with 0 in indegree table into Q
- For each course in Q,
  - dequeue it,
  - Add it into result list
  - decrease 1 in indegree table for all other courses which require it.
  - Add any new courses with 0 indegree.
- Repeat until Q is empty

# LC Graph#4: 417. Pacific Atlantic Water Flow.

```
1 class Solution {
2     const int PAC = 0x1;
3     const int ATL = 0x2;
4
5 public:
6     void dfs(int r, int c, vector<vector<int>>& h, int flag, vector<vector<int>>& map) {
7         vector<vector<int>> DIRS;
8         vector<int> up {-1, 0}, down {1,0}, left{0,-1}, right{0,1};
9         DIRS.push_back(up); DIRS.push_back(down); DIRS.push_back(left); DIRS.push_back(right);
10        int R = h.size();
11        int C = h[0].size();
12        map[r][c] |= flag;
13        // DFS explore thru 4 directions
14        for (auto d: DIRS) {
15            int nr = r + d[0];
16            int nc = c + d[1];
17            if (nr < 0 || nr >= R || nc < 0 || nc >= C) continue; // out of space
18            if ((map[nr][nc] & flag) == flag) continue; // reached already
19            if (h[nr][nc] >= h[r][c]) dfs(nr, nc, h, flag, map); // if neighbor is not lower, explore
20        }
21    }
22    vector<vector<int>> pacificAtlantic(vector<vector<int>>& heights) {
23        int R = heights.size();
24        int C = heights[0].size();
25        vector<vector<int>> map(R, vector<int>(C,0));
26
27        // do DFS traversal starting from known cells:
28        // pacific: top and left
29        // atlantic: bottom and right
30        for (int r = 0; r < R; r++) {
31            dfs(r, 0, heights, PAC, map); // left side
32            dfs(r, C-1, heights, ATL, map); // right side
33        }
34        for (int c = 0; c < C; c++) {
35            dfs(0, c, heights, PAC, map); // top side
36            dfs(R-1, c, heights, ATL, map); // bottom side
37        }
38        vector<vector<int>> res;
39        for (int r = 0; r < R; r++) {
40            for (int c = 0; c < C; c++) {
41                if (map[r][c] == (PAC|ATL)) {
42                    vector<int> ans {r,c};
43                    res.push_back(ans);
44                }
45            }
46        }
47        return res;
48    }
49};
```



- Use DFS find all cells (4 directions) which have equal or higher HEIGHT
- Pacific starts top & left cells; Atlantic starts bottom and right cells
- Use a vector<vector<int>> map to record DFS result
- Check each cell result in map, if both PAC and ATL set, add into result.

# LC Graph#5: 200. Number of Islands.

```
1 class Solution {
2     public:
3         int R, C;
4         void dfs(int r, int c, vector<vector<char>>& grid) {
5             vector<vector<int>> DIRS { {-1,0}, {1,0}, {0,-1}, {0,1}};
6             grid[r][c] = 0; //visit and zero out
7             for (auto d: DIRS) {
8                 int nr = r + d[0];
9                 int nc = c + d[1];
10                if (nr < 0 || nc < 0 || nr >= R || nc >= C) continue;
11                if (grid[nr][nc] == '1') dfs(nr, nc, grid);
12            }
13        }
14
15        int numIslands(vector<vector<char>>& grid) {
16            R = grid.size();
17            C = grid[0].size();
18            int res = 0;
19
20            for (int r = 0; r < R; r++) {
21                for (int c = 0; c < C; c++) {
22                    if (grid[r][c] == '1') {
23                        res++;
24                        dfs(r, c, grid);
25                    }
26                }
27            }
28            return res;
29        }
30    };
```

- Search thru grid, find 1,
- Increment result
- Perform DFS to zero out all neighbors(UP, BOTTOM, LEFT, RIGHT)
- Repeat above



# LC Graph#6: 128. Longest Consecutive Sequence. Must be $O(n)$

```
1 class Solution {  
2 public:  
3     int longestConsecutive(vector<int>& nums) {  
4         if (nums.size() == 0) return 0;  
5         unordered_set<int> allNums(nums.begin(), nums.end());  
6         unordered_map<int, bool> checked; // to ensure  $O(n)$  check  
7         int res = 1;  
8         for (int i = 1; i < nums.size(); i++) {  
9             int n = nums[i];  
10            if (checked.find(n) == checked.end()) {  
11                int counter = 1;  
12                int target = n - 1; // search downwards  
13                while (allNums.find(target) != allNums.end()) {  
14                    counter++;  
15                    checked[target] = true;  
16                    target--;  
17                }  
18                target = n + 1; // search upwards  
19                while (allNums.find(target) != allNums.end()) {  
20                    counter++;  
21                    checked[target] = true;  
22                    target++;  
23                }  
24                res = max(counter, res);  
25                checked[n] = true;  
26            }  
27        }  
28        return res;  
29    }  
30};
```

- Add all numbers into a set
- Go thru each number, search downward and upwards consecutively
- Counter all of them
- Mark them checked in HashMap
- Return max counter

# LC Graph#7: 261. Graph Valid Tree: i.e. all connected without loop.

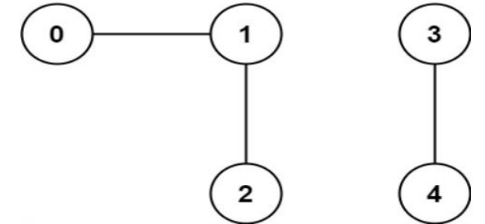
```
1 class Solution {
2 public:
3     bool validTree(int n, vector<vector<int>>& edges) {
4         // establish adjList
5         vector<vector<int>> adj(n);
6         for (int i = 0; i < edges.size(); i++) {
7             int a = edges[i][0]; // use a, b make more sense since this is undirected graph
8             int b = edges[i][1];
9             adj[a].push_back(b); // every edge, we have two edges save, need remove
10            adj[b].push_back(a); // one edge during traverse the first one
11        }
12        unordered_set<int> seen; // save the traversed nodes as result
13        stack<int> st;
14        st.push(0); // start with first node
15        seen.insert(0);
16        while (!st.empty()) {
17            int n = st.top(); st.pop();
18            for (auto nei: adj[n]) {
19                if (seen.find(nei) != seen.end()) {
20                    return false; // there is cycle
21                }
22                seen.insert(nei);
23                st.push(nei);
24
25                // remove edge from nei->n since this is undirected graph
26                vector<int> nodes_2_nei = adj[nei];
27                nodes_2_nei.erase(find(nodes_2_nei.begin(), nodes_2_nei.end(), n));
28                adj[nei] = nodes_2_nei; // need update back!!!
29            }
30        }
31
32        return seen.size() == n;
33    }
34};
```

- A graph is a valid tree: all nodes are connected without loop
- Establish adjacency list for each node by going thru edge list: by directions
- DFS using STACK: Starts with first node, use a “seen” set to track visited node.
- While add each neighbor into stack, remove current node from neighbor’s neighbor list!

# LC Graph#8: 323. Number of Connected Components in an Undirected Graph.

```
1 class Solution {
2     vector<vector<int>>> G;
3     vector<bool> visited;
4 public:
5     void dfs(int i) {
6         visited[i] = true;
7         for (int j = 0; j < G.size(); j++) {
8             if (G[i][j] && visited[j] == false) dfs(j);
9         }
10    }
11
12    int countComponents(int n, vector<vector<int>>& edges) {
13        visited = vector<bool>(n, false);
14        G = vector<vector<int>>(n, vector<int>(n,0));
15        for (int i = 0; i < edges.size(); i++) {
16            G[edges[i][0]][edges[i][1]] = 1;
17            G[edges[i][1]][edges[i][0]] = 1;
18        }
19
20        int res = 0;
21        for (int i = 0; i < n; i++) {
22            if (visited[i] == false) {
23                dfs(i);
24                res++;
25            }
26        }
27
28        return res;
29    }
30};
```

Example 1:



Input: n = 5, edges = [[0,1],[1,2],[3,4]]  
Output: 2

- Transform edges into G
- Use vector<bool> visited(n, false) to track if node is visited by dfs
- Start dfs with node 0, every return of dfs, increment res++;



# LC Graph#9:

## 743: network delay time.

```

1  class Solution {
2  public:
3      int networkDelayTime(vector<vector<int>>& times, int n, int k) {
4          // establish G and adj list for each node
5          int N = n + 1;
6          vector<vector<int>> G(N, vector<int>(N,0));
7          vector<vector<int>> adj(N); // holds neighbors for easy access
8          for (int i = 0; i < times.size(); i++) {
9              int u = times[i][0];
10             int v = times[i][1];
11             int w = times[i][2];
12             G[u][v] = w;
13             vector<int> neis1 = adj[u];
14             neis1.push_back(v);
15             adj[u] = neis1;
16         }
17
18         vector<int> time(n+1, INT_MAX); // hold the final result
19         vector<bool> visited(n+1, false);
20         visited[0] = true; // since we are not using 0 position, exclude it
21         time[k] = 0; // set the given node K as starting point
22

```

```

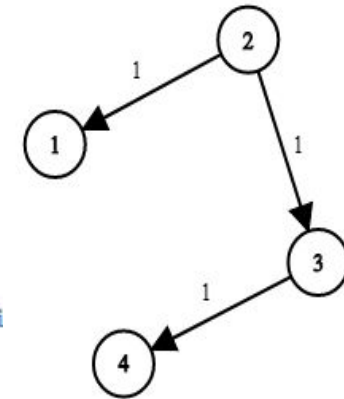
44     }
45     int result = 0;
46     for (int i = 1; i < time.size(); i++)
47         result = max(result, time[i]);
48
49     return result == INT_MAX? -1: result;
50 }
51 };

```

You are given a network of  $n$  nodes, labeled from 1 to  $n$ . You are also given `times`, a list of travel times as directed edges `times[i] = (ui, vi, wi)`, where  $u_i$  is the source node,  $v_i$  is the target node, and  $w_i$  is the time it takes for a signal to travel from source to target.

We will send a signal from a given node  $k$ . Return the time it takes for all the  $n$  nodes to receive the signal. If it is impossible for all the  $n$  nodes to receive the signal, return  $-1$ .

Example 1:



Input: `times = [[2,1,1],[2,3,1],[3,4,1]]`, `n = 4`, `k = 2`  
Output: 2

visited	T	F	T	F	F	F	F	F
time			0			INF	INF	
	k							

- Among NOT visited nodes, find node with mini time as node  $u$
- Find all neighbors of  $u$ , name as  $v$ ,  $\min(\text{time}[v], \text{time}[u] + w)$ , make node  $u$  visited.
- Repeat above until all nodes are visited or  $\text{TIME} = \text{INF}$  (not reachable)

# LC Graph#10: 787: cheapest flights within K stops.

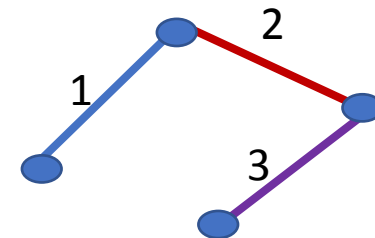
- Key points

# LC# 1135. Connecting Cities With Minimum Cost

## (Prims algorithm)

```
1  class Solution {
2  public:
3      // use prims algorithm:
4      // select the smallest edge, but make sure one node is in node hashset, the other one is not
5      // until N nodes in hashset
6  static bool cmp(vector<int> c1, vector<int> c2) {
7      return c1[2] < c2[2];
8  }
9
10 int minimumCost(int n, vector<vector<int>>& connections) {
11     int res = 0;
12     unordered_set<int> nodes;
13     sort(connections.begin(), connections.end(), cmp); // nLogN
14     //take the first minimum cost edge
15     vector<int> first = connections[0];
16     res = first[2];
17     nodes.insert(first[0]);
18     nodes.insert(first[1]);
19     connections.erase(connections.begin());
20     while (nodes.size() < n) {
21         bool foundConnection = false;
22         for (int i = 0; i < connections.size(); i++) { // N*N
23             vector<int> c = connections[i];
24             bool node1InSet = (nodes.find(c[0]) != nodes.end());
25             bool node2InSet = (nodes.find(c[1]) != nodes.end());
26             if ((node1InSet && !node2InSet) || (!node1InSet && node2InSet)) {
27                 res += c[2];
28                 if (!node1InSet) nodes.insert(c[0]);
29                 if (!node2InSet) nodes.insert(c[1]);
30                 connections.erase(connections.begin()+i);
31                 foundConnection = true;
32                 break;
33             }
34         }
35         if (!foundConnection && (nodes.size() < n)) return -1;
36     }
37     return res;
38 }
39 };
```

- Sort connection first
- Add first mini cost connection, remove it from connections, add nodes into set
- Find mini cost connection with only one node in set until all nodes are in.



Misc: LRU cache, LC#146

# Tree Height, Prime number and Inorder Traversal

```
int hight_dfs(Node* node, int &res) {
    if (node == NULL) return -1;
    if (node->children.size() == 0) return 0;

    vector<int> heights;
    for (auto n: node->children) {
        heights.push_back(hight_dfs(n, res));
    }
    sort(heights.begin(), heights.end(), greater<int>());
    if (heights.size() >= 2) {
        res = max(res, heights[0] + heights[1] + 2); // add two: one for each side
    }
    return heights[0]+1;
}
```

```
class Solution {
public:
    int countPrimes(int n) {
        if (n <= 2) return 0;

        vector<bool> prime_nums(n, true);
        for (int p = 2; p*p < n; p++) {
            if (prime_nums[p]) {
                for (int i=p*p; i < n; i +=p)
                    prime_nums[i] = false;
            }
        }
        int res = 0;
        for (int i = 2; i < n; i++)
            if (prime_nums[i]) res++;

        return res;
    };
};
```

```
vector<int> inorderTraversal(TreeNode* root) {
    vector<int> res;
    if (root == NULL) return res;
#ifdef RECURSIVE
    vector<int> left = preorderTraversal(root->left);
    for (int i = 0; i < left.size(); i++)
        res.push_back(left[i]);
    res.push_back(root->val);
    vector<int> right = preorderTraversal(root->right);
    for (int i = 0; i < right.size(); i++)
        res.push_back(right[i]);
#endif
    TreeNode *cur = root;
    stack<TreeNode *> st;
    while (cur != NULL || !st.empty()) {
        if (cur != NULL) {
            st.push(cur);
            cur = cur->left; // push down all the way to bottom most left node
        } else {
            cur = st.top(); st.pop();
            res.push_back(cur->val);
            cur = cur->right;
        }
    }
    return res;
}
```



# From Excel

- 937s, 408s, 65h,242s, 49m, 1062m, 1092m,14m,12m,527h
- 752m(open lock, BFS), 56s, 80m, 88s, 238s, 16m, 41m, 128m, 35s,
- 69s, 34s, 153m, 162m(peak), 278m, 33m, 81m, 34m, 50m, 96m,
- 762h, 136s, 137m, 169s, 229s, 134m(gas), 179m, 402m, 55m, 45m,
- 31m,19s,21s,82s,83s,86s,206s,92s,61m,109m,
- 138m,141s,142m,143m,148m,538s,110s(balanced tree),102s,107s(Tree Level T),144s (Tree Preorder T),
- 104s,105m(binT from Pre & inoder),297m,98m,285m,510m,366m,156m(BT upDown),39m,207m(course schedule),
- 51m,52m,90m,78m,47m,46m, 63s, 70s, 53s,152m(max product subarray),
- 72m,115h,120m(trangle min sum),139m(word break, recur with memo), 140h (word break II, backtrack), 97m,91m, 128m(int array LongConSeq), 23m,
- 232m, 155m,263m,264m(ugly number), 212h, 79h, 295h,84h, 438s,311m,288m,
- 981m(time based key store), 706m(Design HashMap)

# Java: HashMap and HashSet

```
HashMap<Integer, Integer> map = new HashMap<Integer, Integer>();  
  
/* Add new entry */  
map.put(1, 100); map.put(2, 200); if (map.size() == 2) System.out.println("Cool");  
  
/* Get/check using key */  
Integer val = map.get(1);  
  
/* Loop thru key */ for (Integer key: map.keySet()) System.out.println("key is " + key.intValue());  
/* Loop thru value */  
for (Integer val: map.values()) System.out.println("val is " + val.intValue());  
  
/* Iterator using Map.Entry*/  
for (Map.Entry<Integer, Integer> ent: map.entrySet())  
    System.out.println("Key is " + ent.getKey() + ", value is " + ent.getValue());  
  
HashSet<String> set = new HashSet<String>();  
set.add("xyz"); if (set.contains("abc")) System.out.println("abc is not found");
```

# Java: Stack and Queue

```
Stack<Integer> st = new Stack<Integer>;  
st.push(1); st.push(2); st.pop(); st.peek(); st.size();  
while (!st.isEmpty()) System.out.println("pop: " + st.pop());
```

```
Queue<Integer> q = new LinkedList<Integer>();  
q.add(1); q.add(2); q.add(3);  
System.out.println("Q size is " + q.size());  
while (!q.isEmpty()) System.out.println("Q remove : " + q.remove()); // remove() throw exp if empty  
Integer val = q.poll(); //remove() method returns the head of the queue and removes it.  
                // It returns null if the queue is empty.
```

```
Vector<Integer> v = new Vector<Integer>(); // thread safe  
v.add(1); v.add(2); Integer val = v.get(1);  
v.remove(0); v.remove(new Integer(2));  
for (int i = 0; i < v.size(); i++) System.out.print(v.get(i) + " ");
```

```
ArrayList<Integer> al = new ArrayList <Integer>(); // not synchronized, faster  
al.add(1); al.add(2); Integer val = al.get(1);  
al.remove(0); al.remove(new Integer(2));  
for (int i = 0; i < al.size(); i++) System.out.print(al.get(i) + " ");
```

# Java: util.Arrays

- `Java.util.Arrays.binarySearch(int[] arr, int key):`
- `Java.util.Arrays.copyOf(int[] arr, int newLen)`
- `Java.util.Arrays.sort(int[] arr)` or  
`Java.util.Arrays.sort(int[] arr, Collections.reverseOrder())`
- Class `MyCmp` implements `Comparator<Student>` {  
    `public int compare(Student a, Student b) { return a.id – b.id;} }`
- `Java.util.Arrays.sort(Student[] arr, new MyCmp())`

# Python: list (just like vector in C++ std)

- `L1 = [1,2,3]`   `L2 = list()`
- `L1.append(4)`   `L2.append("abc")`
- `L1.insert(0, 100)` # add 100 at the beginning (index, val)
- `val1 = L1.pop(0)` # remove first element and return
- `Val2 = L1.pop()` # pop last element, for this reason, list can be used as stack
- `L1.remove(2)` # remove the first occurrence of value 2
- `L1.reverse()`
- `L1.sort()`
- `L1.count(2)` // tell how many occurrence of value 2
- `L1.index(3)`   `L1[3]` // return value at given index
- `Size = len(L1)` # return how many elements in the list



# Python dict: (unordered\_map in C++ std)

- `D1 = { 1:"abc", 2:"xyz"} D2 = dict()`
- `D1[100] = "odfag" D2["xyz"] = 987 // add new element`
- `len(D1)` # get the total elements in dictionary
- `del D1[2]` #remove element using key
- `D1.pop(2)` # another way to remove element using key
- `2 in D1:` # tell if key exists
- `"abc" in D1.values():` # tell if value exists
- `for k in D1: print(k); print(D1[k])` # iterate thru all elements

# Python queue & stack

```
from queue import Queue
q = Queue(maxsize = 3)
q.put('a'); q.put('b'); q.put('c')
print(q.get())
q.empty()
q.full()
q.qsize()
```

```
from queue import LifoQueue
st = LifoQueue(maxsize = 3)
st.put('a'); st.put('b'); st.put('c')
print(st.get())
st.empty()
st.full()
st.qsize()
```

